Data Structures

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Fall, 2000

CHAPTER 20
Hash Tables

Contents

We will see:

- Several methods of implementing the hash tables;
- Comparisons of these methods;
- Some applications of hashing; and
- Comparisons of hash tables and binary search trees.

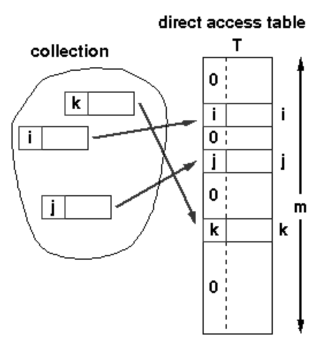
Hash Tables

- All search structures so far
  - Relied on a comparison operation
  - Performance $O(n)$ or $O(\log n)$
- Assume I have a function
  - $f(\ key\ ) \Rightarrow integer$

    ie one that maps a key to an integer
- What performance might I expect now?

Hash Tables - Structure

- Simplest case:
  - Assume items have integer keys in the range $1 .. m$
  - Use the value of the key itself to select a slot in a direct access table in which to store the item
  - To search for an item with key, k, just look in slot k
    - If there's an item there, you've found it
    - If the tag is 0, it's missing.
  - Constant time, $O(1)$

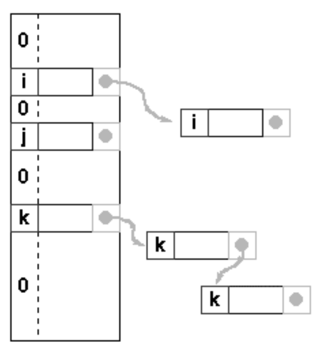


Hash Tables - Constraints

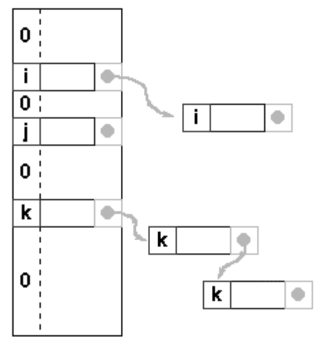
- Constraints
  - Keys must be unique
  - Keys must lie in a small range
  - For storage efficiency,
    keys must be dense in the range
  - If they're sparse (lots of gaps between values),
    a lot of space is used to obtain speed
    - Space for speed trade-off

Hash Tables - Relaxing the constraints

- Keys must be unique
  - Construct a linked list of duplicates "attached" to each slot
  - If a search can be satisfied by *any* item with key, k, performance is still $O(1)$
    but
  - If the item has some other distinguishing feature which must be matched, we get $O(n^{max})$
    where n^{max} is the largest number of duplicates - or length of the longest chain



Hash Tables - Relaxing the constraints



- Keys are integers
  - Need a hash function
    h(key) ➔ *integer*
    ie one that maps a key to an integer
  - Applying this function to the key produces an address
  - If h maps each key to a *unique integer* in the range $0 .. m-1$ then search is $O(1)$

Hash Tables - Hash functions

- Form of the hash function
  - Example - using an *n*-character key

```
int hash( char *s, int n ) {
    int sum = 0;
    while( n-- ) sum = sum + *s++;
    return sum % 256;
    }
```

    returns a value in 0 .. 255
  - `xor` function is also commonly used

```
sum = sum ^ *s++;
```

  - But any function that generates integers in $0..m-1$ for some suitable (*not too large*) *m* will do
  - As long as the hash function itself is $O(1)$!

Hash Tables - Collisions

- Hash function
  - With this hash function

```
int hash( char *s, int n ) {
    int sum = 0;
    while( n-- ) sum = sum + *s++;
    return sum % 256;
    }
```

  - `hash( "AB", 2 )` and `hash( "BA", 2 )` return the same value!
  - This is called a collision
  - A variety of techniques are used for resolving collisions

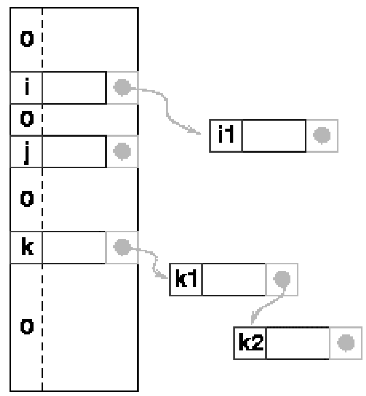
Hash Tables - Collision handling

- Collisions
  - Occur when the hash function maps two different keys to the same address
  - The table must be able to recognize and resolve this
  - Recognize
    - Store the actual key with the item in the hash table
    - Compute the address
      - k = h(key)
    - Check for a hit
      - *if (* table[k].key == key *) then hit*
        else try next entry
  - Resolution
    - Variety of techniques

We'll look at various "*try next entry*" schemes

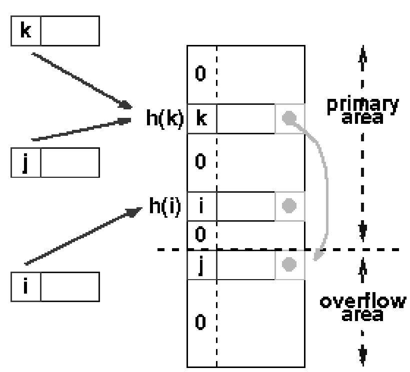
Hash Tables - Linked lists

- Collisions - Resolution
  1. Linked list attached to each primary table slot
     - h(i) == h(i1)
     - h(k) == h(k1) == h(k2)
  - Searching for i1
     - Calculate h(i1)
     - Item in table, i, doesn't match
     - Follow linked list to i1
  - If `NULL` found, key isn't in table



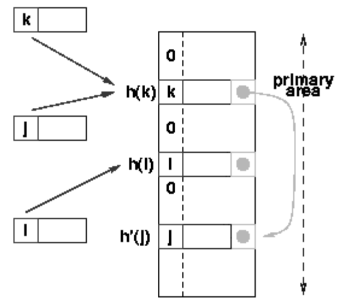
Hash Tables - Overflow area

2. Overflow area
   - Linked list constructed in special area of table called overflow area
- h(k) == h(j)
- k stored first
- Adding j
  - Calculate h(j)
  - Find k
  - Get first slot in overflow area
  - Put j in it
  - k's pointer points to this slot
- Searching - same as linked list



Hash Tables - Re-hashing

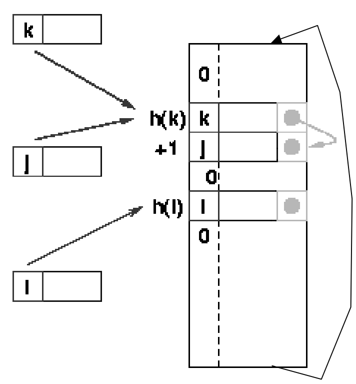
3. Use a second hash function
   - Many variations
   - General term: re-hashing
- h(k) == h(j)
- k stored first
- Adding j
  - Calculate h(j)
  - Find k
  - Repeat until we find an empty slot
    - Calculate h'(j)
  - Put j in it
- Searching - Use h(x), then h'(x)



h'(x) -
second hash function

Hash Tables - Re-hash functions

3. The re-hash function
   - Many variations
- Linear probing
  - h'(x) is +1
  - Go to the next slot until you find one empty



- Can lead to bad clustering
- Re-hash keys fill in gaps between other keys and exacerbate the collision problem

Hash Tables - Re-hash functions

3. The re-hash function
   - Many variations
- Quadratic probing
   - $h'(x)$ is $c\,i^2$ on the i^{th} probe
   - Avoids primary clustering
   - Secondary clustering occurs
     - All keys which collide on $h(x)$ follow the same sequence
     - First
       $a = h(j) = h(k)$
     - Then $a + c$, $a + 4c$, $a + 9c$, ...
     - Secondary clustering generally less of a problem

Hash Tables - Collision Resolution Summary

- Chaining
  + Unlimited number of elements
  + Unlimited number of collisions
  - Overhead of multiple linked lists
- Re-hashing
  + Fast re-hashing
  + Fast access through use of main table space
  - Maximum number of elements must be known
  - Multiple collisions become probable
- Overflow area
  + Fast access
  + Collisions don't use primary table space
  - Two parameters which govern performance need to be estimated

Hash Tables - Summary so far ...

- Potential $O(1)$ search time
  - If a suitable function *h(key)* → *integer* can be found
- Space for speed trade-off
  - "Full" hash tables don't work (more later!)
- Collisions
  - Inevitable
    - Hash function reduces amount of information in key
  - Various resolution strategies
    - Linked lists
    - Overflow areas
    - Re-hash functions
      - Linear probing h' is $+1$
      - Quadratic probing h' is $+ci^2$
      - Any other hash function!
        or even sequence of functions!

Hash Tables - Choosing the Hash Function

- "Almost any function will do"
  - But some functions are definitely better than others!
- Key criterion
  - Minimum number of collisions
    - Keeps chains short
    - Maintains $O(1)$ average

Hash Tables - Choosing the Hash Function

- Uniform hashing
  - Ideal hash function: the number of keys that map to each slot should be equal
    - $P(k)$ = probability that a key, k, occurs
    - If there are m slots in our hash table,
    - a uniform hashing function, $h(k)$, would ensure:

$$\sum_{k \mid h(k)=0} P(k) = \sum_{k \mid h(k)=1} P(k) = \ldots\ldots \sum_{k \mid h(k)=m-1} P(k) = \frac{1}{m}$$

Read as sum over all k such that $h(k) = 0$

- *or, in plain English,*
- the number of keys that map to each slot is equal

Hash Tables - A Uniform Hash Function

- *If* the keys are integers randomly distributed in $[\,0\,,\,r\,)$, **Read as $0 \le k < r$**

  then

$$h(k) = \left\lfloor \frac{mk}{r} \right\rfloor$$

  is a uniform hash function
- Most hashing functions can be made to map the keys to $[\,0\,,\,r\,)$ for some r
  - *eg* adding the ASCII codes for characters mod 255 will give values in [0, 256) or [0, 255]
  - Replace + by xor
    - ➧ same range without the mod operation

Hash Tables - Reducing the range to [0, m)

- We've mapped the keys to a range of integers
  $0 \leq k < r$
- Now we must reduce this range to [0, m)
  where m is a reasonable size for the hash table

- Strategies
  1. Division - use a mod function
  2. Multiplication
  3. Universal hashing

Hash Tables - Reducing the range to [0, m)

1. Division
   - Use a mod function

     $h(k) = k \bmod m$

   – Choice of m?
     - Powers of 2 are generally not good!

       $h(k) = k \bmod 2^n$

       selects last n bits of k

       k mod 2^8 selects these bits

       `01100101110` $\boxed{\texttt{00011010}}$

       – All combinations are not generally equally likely

   – Prime numbers close to 2^n seem to be good choices

     eg want ~4000 entry table, choose m = 4093

Hash Tables - Reducing the range to [0, m)

2. Multiplication method
   - Multiply the key by constant, A, 0 < A < 1
   - Extract the fractional part of the product

$$(kA - \lfloor kA \rfloor)$$

   - Multiply this by m

$$h(k) = \lfloor m * (kA - \lfloor kA \rfloor) \rfloor$$

   – Now m is not critical and a power of 2 can be chosen
   – So this procedure is fast on a typical digital computer
      - A = ½($\sqrt{5}$ -1) seems to be a good choice (*see* Knuth)

Hash Tables - Reducing the range to [0, m)

3. Universal Hashing
   - A determined "adversary" can always find a set of data that will defeat any hash function
      - Hash all keys to same slot ➔ $O(n)$ search
   – Select the hash function randomly (at run time) from a set of hash functions

   ➔Reduced probability of poor performance
   – Set of functions, $\mathcal{H}$, which map keys to [0, m)
   – $\mathcal{H}$, is universal, if for each pair of keys, x and y, the number of functions, h from $\mathcal{H}$, for which $h(x) = h(y)$ is $|\mathcal{H}|/m$

Hash Tables - Reducing the range to [0, m)

3. Universal Hashing

- Can we design a set of universal hash functions?
- *Quite easily*
- Key, $x = x_0, x_1, x_2, \ldots, x_r$
- Choose $a = \langle a_0, a_1, a_2, \ldots, a_r \rangle$
  a is a sequence of elements
  chosen randomly from $\{ 0, m\text{-}1 \}$
- $h_a(x) = \Sigma\, a_i x_i \bmod m$
- There are m^{r+1} *sequences a,*
  so there are m^{r+1} functions, $h_a(x)$
- Theorem
  - The h_a form a set of universal hash functions

x_0	x_1	x_2	....	x_r

n*-bit "bytes" of *x

Proof:
***See* Cormen …**

Collision Frequency

- Birthdays *or* the von Mises paradox
  - There are 365 days in a normal year
    Birthdays on the same day unlikely?
  - How many people do I need before "it's an even bet" (*ie* the probability is > 50%) that two have the same birthday?
  - View
    - the days of the year as the slots in a hash table
    - the "birthday function" as mapping people to slots
  - Answering von Mises' question answers the question about the probability of collisions in a hash table

Distinct Birthdays

- Let *Q(n)* = probability that *n* people have distinct birthdays
- *Q*(1) = 1
- With two people, the 2nd has only 364 "free" birthdays

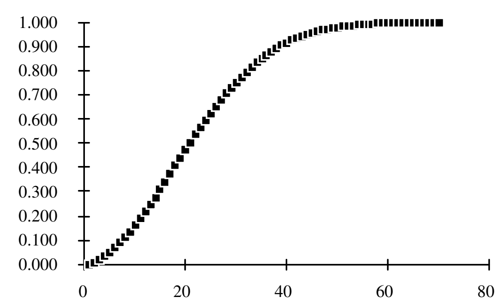
$$Q(2) = Q(1) * \frac{364}{365}$$

- The 3rd has only 363, and so on:

$$Q(n) = Q(1) * \frac{364}{365} * \frac{363}{365} * \ldots * \frac{365-n+1}{365}$$

Coincident Birthdays

- Probability of having two identical birthdays
- *P(n)* = 1 - *Q(n)*
- *P*(23) = 0.507
- With 23 entries, table is only 23/365 = 6.3% full!



Hash Tables - Load factor

- Collisions are very probable!
- Table load factor

$$\alpha = \frac{n}{m}$$

n = number of items
m = number of slots

must be kept low

- Detailed analyses of the average chain length (or number of comparisons/search) are available
- Separate chaining
  - linked lists attached to each slot

gives best performance

  - but uses more space!

Hash Tables - General Design

1. Choose the table size
   - Large tables reduce the probability of collisions!
   - Table size, m
   - n items
   - Collision probability $\alpha = n / m$
2. Choose a table organization
   - Does the collection keep growing?
     - Linked lists (....... but consider a tree!)
   - Size relatively static?
     - Overflow area *or*
     - Re-hash

Hash Tables - General Design

3. Choose a hash function
   - A simple (and fast) one may well be fine ...
   - Read your text for some ideas!
4. Check the hash function against your data

   Fixed data
   - Try various h, m
     until the maximum collision chain is acceptable

   ➔Known performance

   Changing data
   - Choose some representative data
   - Try various h, m until collision chain is OK

   ➔Usually predictable performance

Hash Tables - Review

- *If you can meet the constraints*
  + Hash Tables will generally give good performance
  + $O(1)$ search
- Use a hash table instead of a binary search tree if you do not need order statistics and are worried about nonrandom inputs.
- Hashing Applications:
  1. Symbol tables (compilers use hash tables to keep track of declared variables in source code)
  2. Transposition table (in game programs)
  3. In online spelling checkers