## Implementation

## Objectives

- Introduce basic implementation strategies
- Clipping
- Scan conversion
- Introduce clipping algorithms for polygons
- Survey hidden-surface algorithms

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## Meta Algorithms

- Consider two approaches to rendering a scene with opaque objects
- For every pixel, determine which object that projects on the pixel is closest to the viewer and compute the shade of this pixel
- Ray tracing paradigm
- For every object, determine which pixels it covers and shade these pixels
- Pipeline approach
- Must keep track of depths

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| :--- |
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## Implementation (ctd)

## Objectives

- Survey Line Drawing Algorithms
- DDA
- Bresenham

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## Common Tasks

- Clipping
- Rasterization or scan conversion
- Antialiasing
- Transformations
- Hidden surface removal


[^0]
## Clipping

- 2D against clipping window
- 3D against clipping volume
- Easy for line segments polygons
- Hard for curves and text
- Convert to lines and polygons first


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## Cohen-Sutherland Algorithm

- Idea: eliminate as many cases as possible without computing intersections
- Start with four lines that determine the sides of the clipping window


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## Clipping 2D Line Segments

- Brute force approach: compute intersections with all sides of clipping window
- Inefficient: one division per intersection


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## four lines

- Draw (accept) line segment as is

- Case 2: both endpoints outside all lines and on same side of some line
- Discard (reject) the line segment

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## The Cases

- Case 3: One endpoint inside, one outside
- Must do at least one intersection
- Case 4: Both outside all lines but not on same side of any line
- May have part inside
- Must do af least one intersection




## Defining Outcodes

- For each endpoint, define an outcode

| $\mathrm{b}_{0} \mathrm{~b}_{1} \mathrm{~b}_{2} \mathrm{~b}_{3}$ | 1001 | 1000 | 1010 |
| :---: | :---: | :---: | :---: |
| $\mathrm{b}_{0}=1$ if $\mathrm{y}>\mathrm{y}_{\text {max }}, 0$ otherwise | 0001 | 0000 | 0010 |
| $\mathrm{b}_{1}=1$ if $\mathrm{y}<\mathrm{y}_{\text {min }}, 0$ otherwise | 0101 | 0100 | 0110 |
| $\mathrm{b}_{2}=1$ if $\mathrm{x}>\mathrm{x}_{\text {max }}, 0$ otherwise | 0101 | 0100 | 010 |

- Outcodes divide space into 9 regions
- Computation of outcode requires at most 4 subtractions

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## Using Outcodes

- CD: outcode $(C)=0$, outcode $(D)=0010 \neq 0$
- Compute intersection
- Location of 1 in outcode(D) determines which edge to intersect with
- Note if there were a segment from A to a point in a region with 2 ones in outcode, we might have to do two interesections


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## Using Outcodes

-EF: outcode(E) logically ANDed with outcode(F) (bitwise) $\neq 0$

- Both outcodes have a 1 bit in the same place
- Line segment is outside of corresponding side of clipping window
- reject



## Efficiency

- In many applications, the clipping window is small relative to the size of the whole data base
- Most line segments are outside one or more side of the window and can be eliminated based on their outcodes
- Inefficiency when code has to be reexecuted for line segments that must be shortened in more than one step

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## Using Outcodes

- GH and IJ: same outcodes, neither zero but logical AND yields zero
- Shorten line segment by intersecting with one of sides of window
- Compute outcode of intersection (new endpoint of shortened line segment)
- Reexecute algorithm

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## Cohen Sutherland in 3D

- Use 6-bit outcodes
- When needed, clip line segment against planes


[^1]
## Liang-Barsky Clipping

- Consider the parametric form of a line segment

$$
\mathbf{p}(\alpha)=(1-\alpha) \mathbf{p}_{1}+\alpha \mathbf{p}_{2} \quad 1 \geq \alpha \geq 0
$$



- We can distinguish between the cases by looking at the ordering of the values of $\alpha$ where the line determined by the line segment crosses the lines that determine the window

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## Advantages

- Can accept/reject as easily as with Cohen-Sutherland
- Decisions can be made without calculating intersections (without divisions)
- Using values of $\alpha$, we do not have to use algorithm recursively as with C-S
- Extends to 3D


## Liang-Barsky Clipping

- In (a): $\alpha_{4}>\alpha_{3}>\alpha_{2}>\alpha_{1}$
- Intersect right, top, left, bottom: shorten
- In (b): $\alpha_{4}>\alpha_{2}>\alpha_{3}>\alpha_{1}$
- Intersect right, left, top, bottom: reject

(a)
(b)

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## Clipping and Normalization

- General clipping in 3D requires intersection of line segments against arbitrary plane
- Example: oblique view


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## Plane-Line Intersections

- Intersection requires 6 multiplications and a division


$$
a=\frac{n \bullet\left(p_{o}-p_{1}\right)}{n \bullet\left(p_{2}-p_{1}\right)}
$$

## Polygon Clipping

- Not as simple as line segment clipping
- Clipping a line segment yields at most one line segment
- Clipping a polygon can yield multiple polygons

- However, clipping a convex polygon can yield at most one other polygon

[^2]KENT STATE ${ }^{23}$

## Normalized Form

top view

before normalization

after normalization
Normalization is part of viewing (pre clipping) but after normalization, we clip against sides of right parallelepiped

Typical intersection calculation now requires only a floating point subtraction, e.g. is $\mathrm{x}>\mathrm{x}_{\text {max }}$ ?
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## Tessellation and Convexity

- One strategy is to replace nonconvex (concave) polygons with a set of triangular polygons (a tessellation)
- Also makes fill easier
- Tessellation code in GLU library


[^3]
## Clipping as a Black Box

- Can consider line segment clipping as a process that takes in two vertices and produces either no vertices or the vertices of a clipped line segment



## Pipeline Clipping of Polygons



- Three dimensions: add front and back clippers
- Strategy used in SGI Geometry Engine
- Small increase in latency


## Pipeline Clipping of Line

 Segments- Clipping against each side of window is independent of other sides
- Can use four independent clippers in a pipeline


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## Bounding Boxes

- Rather than doing clipping on a complex polygon, we can use an axis-aligned bounding box or extent
- Smallest rectangle aligned with axes that encloses the polygon
- Simple to compute: max and $\min$ of $x$ and $y$


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## Bounding boxes

Can usually determine accept/reject based only on bounding box


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Hidden Surface Removal

- Object-space approach: use pairwise testing between polygons (objects)

partially obscuring

can draw independently
- Worst case complexity $\mathrm{O}\left(\mathrm{n}^{2}\right)$ for n polygons

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## Clipping and Visibility

- Clipping has much in common with hidden-surface removal
- In both cases, we are trying to remove objects that are not visible to the camera
- Often we can use visibility or occlusion testing early in the process to eliminate as many polygons as possible before going through the entire pipeline

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## Painter's Algorithm

- Render polygons in back to front order so that polygons behind others are simply painted over

$B$ behind $A$ as seen by viewer


Fill $B$ then $A$

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## Depth Sort

- Requires ordering of polygons first
- O(n $\log n$ ) calculation for ordering
- Not every polygon is either in front or behind all other polygons
- Order polygons and deal with easy cases first, harder later

Polygons sorted by distance from COP ter

$$
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$$



## Easy Cases

- A lies behind all other polygons
- Can render
- Polygons overlap in z but not in either x or y
- Can render independently


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## Back-Face Removal (Culling)

-face is visible iff $90 \geq \theta \geq-90$ equivalently $\cos \theta \geq 0$
or $\mathbf{v} \cdot \mathbf{n} \geq 0$

-plane of face has form ax + by $+\mathrm{cz}+\mathrm{d}=0$ but after normalization $\mathbf{n}=\left(\begin{array}{lll}0 & 0 & 1\end{array}\right)^{\mathrm{T}}$
-need only test the sign of c
-In OpenGL we can simply enable culling
but may not work correctly if we have nonconvex objects
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## Image Space Approach

- Look at each projector (nm for an nx m frame buffer) and find closest of $k$ polygons
- Complexity O(nmk)
- Variations
- Ray tracing
- z-buffer


## Efficiency

- If we work scan line by scan line as we move across a scan line, the depth changes satisfy $a \Delta x+b \Delta y+c \Delta z=0$

$$
\begin{aligned}
& \text { Along scan line } \\
& \Delta \mathrm{y}=0 \\
& \Delta \mathrm{z}=-\frac{a}{c} \Delta \mathrm{x}
\end{aligned}
$$

In screen space $\Delta \mathrm{x}=1$
So only need $\frac{a}{c}$ (a const)


[^5]KENT STATE ${ }^{39}$

## Z-Buffer Algorithm

- Use a buffer called the z or depth buffer to store the depth of the closest object at each pixel found so far
- As we render each polygon, compare the depth of each pixel to depth in z buffer
- If less, place shade of pixel in color buffer and update $z$ buffer

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## Scan-Line Algorithm

- Can combine shading and hidden surface removal through scan line algorithm

scan line i: no need for depth information, can only be in no or one polygon
scan line j: need depth information only when in more than one polygon

[^6]
## Implementation

- Need a data structure to store
- Flag for each polygon (inside/outside)
- Incremental structure for scan lines that stores which edges are encountered
- Parameters for planes



## Visibility Testing

- In many realtime applications, such as games, we want to eliminate as many objects as possible within the application
- Reduce burden on pipeline
- Reduce traffic on bus
-Partition space with Binary Spatial Partition (BSP) Tree

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## BSP Tree

- Can continue recursively
- Plane of C separates B from A
- Plane of D separates E and F
- Can put this information in a BSP tree
- Use for visibility and occlusion testing


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## Rasterization

## Rasterization or Scan Conversion of Lines

-Rasterization (scan conversion)

- Shade pixels that are inside object specified by a set of vertices
- Line segments
- Polygons: scan conversion = fill
- Shades determined by color, texture, shading model
- Here we study algorithms for determining the correct pixels starting with the vertices


## Line Drawing - Algorithm 1

```
A Straightforward Implementation
Drawline(x1,y1,x2,y2)
    int x1,y1,x2,y2;
    {
        float y;
    int x;
    for (x=x1; x<=x2; x++) {
    y=y1 + (x-x1)*(y2-y1)/(x2-x1)
        SetPixel(x, Round(y) );
    }
}
- Such a line should ideally have the following properties.
- Straight,
- pass through endpoints
- smooth
- independent of endpoint order
- uniform brightness
- brightness independent of slope
- efficient

\section*{Scan Conversion of Line Segments}
- Start with line segment in window coordinates with integer values for endpoints
- Assume implementation has a write_pixel function
\[
m=\frac{\Delta y}{\Delta x}
\]


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}

\section*{DDA Algorithm}
- Digital Differential Analyzer
- DDA was a mechanical device for numerical solution of differential equations
- Line \(y=m x+h\) satisfies differential equation \(d y / d x=m=\Delta y / \Delta x=y_{2}-y_{1} / x_{2}-x_{1}\)
- Along scan line \(\Delta x=1\)
```

For(x=x1; x<=x2, x++) {
y+=m;
write_pixel(x, round(y), line_color)
}

```

\section*{Line Drawing Algorithm Comparison}
- Advantages over Algorithm 1
- eliminates multiplication
- improves speed
- Disadvantages
- round-off error builds up
- get pixel drift
- rounding and fp arithmetic still time consuming
- works well only for \(|\mathrm{m}|<1\)
- need to loop in y for \(|m|>1\)
- need to handle special cases

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\section*{Line Drawing - Algorithm 2}

A Better Implementation
DrawLine( \(\mathrm{x} 1, \mathrm{y} 1, \mathrm{x} 2, \mathrm{y} 2\) )
int \(\mathrm{x} 1, \mathrm{y} 1, \mathrm{x} 2, \mathrm{y} 2\);
\{
float \(m, y\);
int dx,dy, x;
\(\mathrm{dx}=\mathrm{x} 2-\mathrm{x} 1\);
\(\mathrm{dy}=\mathrm{y} 2-\mathrm{y} 1\);
\(\mathrm{m}=\mathrm{dy} / \mathrm{dx}\);
\(y=y 1+0.5 ;\)
for ( \(x=x 1 ; x<=x 2 ; x++\) ) \{
SetPixel(x, Floor(y) );
\(y=y+m ;\)
\}

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\section*{Problem}
- DDA = for each \(x\) plot pixel at closest \(y\)
- Problems for steep lines


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\section*{Using Symmetry}
- Use for \(1 \geq m \geq 0\)
- For \(m>1\), swap role of \(x\) and \(y\)
- For each \(y\), plot closest \(x\)


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\section*{Based on Implicit \\ Representation}
- Explicit: \(y=f(x)\)
- \(y=m(x-x 0)+y 0\) where \(m=d y / d x\)
- Implicit: \(f(x, y)=0\)
\(-F(x, y)=(x-x 0) d y-(y-y 0) d x\)
- if \(F(x, y)=0\) then \((x, y)\) is on line
- \(F(x, y)>0\) then \((x, y)\) is below line
- \(F(x, y)<0\) then \((x, y)\) is above line

\section*{Line Drawing - Midpoint Algorithm}
- The Midpoint or Bresenham's Algorithm
- Uses only integer calculations. It treats line drawing as a sequence of decisions. For each pixel that is drawn the next pixel will be either E or NE , as shown below.


\section*{Midpoint Algorithm}
- The midpoint algorithm makes use of the the implicit definition of the line, \(F(x, y)=0\). The N/NE decisions are made as follows.
- \(\mathrm{d}=\mathrm{F}\left(\mathrm{x}_{\mathrm{p}}+1, \mathrm{y}_{\mathrm{p}}+0.5\right)\)
- if \(d<0\) line below midpt choose \(E\)
- if \(d>0\) line above midpt choose NE
- if \(E\) is chosen

- \(d_{\text {new }}=F\left(x_{p}+2, y_{p}+0.5\right)\)
- \(d_{\text {new }}-d_{\text {old }}=F\left(x_{p}+2, y_{p}+0.5\right)-\)
- \(F\left(x_{p}+1, y_{p}+0.5\right)\)
- \(\Delta d=d_{\text {new }}-d_{\text {old }}=d y\)

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\section*{Midpoint Algorithm}
- If \(N E\) is chosen
- \(d_{\text {new }}=F\left(x_{p}+2, y_{p}+1.5\right)\)
- \(\Delta \mathrm{d}=\mathrm{dy}\) - dx
- Initialization
\(-d_{\text {start }}=F\left(x_{0}+1, y_{0}+0.5\right)=\left(x_{0}+\right.\)
\(\left.1-x_{0}\right) d y-\left(y_{0}+0.5-y_{0}\right) d x\)
\(\quad=d y-d x / 2\)
- Integer only algorithm
\[
-F^{\prime}(x, y)=2 F(x, y) \quad ; d^{\prime}=2 d
\]
- \(d_{\text {start }}^{\prime}=2 d y-d x\)
- \(\Delta d^{\prime}=2 \Delta d\)

\section*{General Bresenham's Algorithm}
- To generalize to lines with arbitrary slope
- consider symmetry between various octants and quadrants
- for \(m>1\), interchange roles of \(x\) and \(y\), that is step in \(y\) direction, and decide whether \(x\) value is above or below line
- if \(m>1\), and right endpoint is the first point, both \(x\) and \(y\) decrease. To ensure uniqueness, independent of direction, always choose upper (or lower) point if the line goes through the mid-point
- handle special cases without invoking algorithm: horizontal, vertical and diagonal lines
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Midpoint Algorithm for x1 < x2 and slope <= 1
diavilime(x1, y1, \(x 2\), yz, colt
int \(x 1, y 1, \times 2, y 2\), colour;
\{
int \(d x, d y, d, i n c E, \operatorname{incNE}, x, y\)
\(\mathrm{dx}=\mathrm{x} 2-\mathrm{x} 1\);
\(\mathrm{dy}=\mathrm{y} 2-\mathrm{y} 1\);
\(d=2^{*} d y-d x\);
incE \(=2^{*}\) dy;
incNE \(=2^{\star}(d y-d x)\);
\(\mathrm{y}=\mathrm{y} 1\);
for \((x=x 1 ; x<=x 2 ; x++)\{\)
setpixel( \(x, y\), colour);
if \((\mathrm{d}>0)\) \{
\(\mathrm{d}=\mathrm{d}+\mathrm{incNE} ;\)
\(y=y+1 ;\)
\} else \{
\(d=d+i n c E ;\)
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\section*{Additional Issues}
- End-point order
- cannot just interchange end-points
- does not work when we use line styles since we need the pattern to go the same way on all segments of a polygon
- varying the intensity of a line with the slope
- consider horizontal line and diagonal line
- both have same number of pixels
- diagonal \(\sqrt{ } 2\) times horizontal line in length
- intensity per unit length less for diagonal

\section*{Polygon Scan Conversion}
- Scan Conversion = Fill
- How to tell inside from outside
- Convex easy
- Nonsimple difficult
- Odd even test
- Count edge crossings
- Winding number

\section*{Winding Number}
- Count clockwise encirclements of point
winding number = 1

- Alternate definition of inside: inside if winding number \(\neq 0\)

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\section*{Filling in the Frame Buffer}
- Fill at end of pipeline
- Convex Polygons only
- Nonconvex polygons assumed to have been tessellated
- Shades (colors) have been computed for vertices (Gouraud shading)
- Combine with z-buffer algorithm
- March across scan lines interpolating shades
- Incremental work small

\section*{Using Interpolation}
\(\mathrm{C}_{1} \mathrm{C}_{2} \mathrm{C}_{3}\) specified by glColor or by vertex shading \(\mathrm{C}_{4}\) determined by interpolating between \(\mathrm{C}_{1}\) and \(\mathrm{C}_{2}\) \(\mathrm{C}_{5}\) determined by interpolating between \(\mathrm{C}_{2}\) and \(\mathrm{C}_{3}\) interpolate between \(\mathrm{C}_{4}\) and \(\mathrm{C}_{5}\) along span


\section*{Flood Fill}
- Fill can be done recursively if we know a seed point located inside, currently background color (WHITE)
- Scan convert edges into buffer in edge/fill color (BLACK)
```

flood_fill(int x, int y) {
if(read_pixel(x,y)= = WHITE) {
write_pixel(x,y,BLACK);
flood_fill(x-1, y);
flood_fill(x+1, y);
flood_fill(x, y+1);
flood_fill(x, y-1);
} }

```
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\section*{Scan Conversion Algorithm}
- intersect each scan-line with all edges
- sort intersections by increasing \(x\) coordinate
- calculate parity of intersections to
 determine in/out
- parity starts even - each intersection inverts
- fill the 'in' pixels - those with odd parity
- General issues - how to handle

Anintersection atolintegerisamelifnactionalent STATE \({ }^{67}\)

\section*{Scan Line Fill}
- Can also fill by maintaining a data structure of all intersections of polygons with scan lines
- Sort by scan line
- Fill each span

vertex order generated by vertex list

desired order

\section*{Scan Conversion Algorithm}
- General issues - how to handle intersection at integer and fractional \(x\) values
- Special cases:
- shared vertices lying on scan-lines double intersections
- count \(y_{\text {min }}\) vertices but not \(y_{\text {max }}\) vertices in parity count
- do NOT count vertices of horizontal edges

\section*{Fractional and Integer} Intersections
- Fractional intersections
- if approaching intersection to the right to determine inside pixel
- take floor if inside, ceil if outside
- Integer intersections
- if leftmost pixel
- make interior,
- rightmost exterior


\section*{Using Spatial Coherence}
- Efficiency can be improved by using spatial coherence
- Edges that intersect scan-line i are likely to intersect i+1
- \(x_{i}\) changes predictably from scan-line \(i\) to i+1
- use an incremental algorithm that calculates the scan-line extrema from extrema of previous scan line by using
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\section*{Aliasing}
- Ideal rasterized line should be 1 pixel wide

- Choosing best y for each x (or visa versa) produces aliased raster lines

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\section*{Antialiasing by Area Averaging}
- Color multiple pixels for each \(x\) depending on coverage by ideal line


\section*{Unweighted Area Sampling}
- Assume background white - lines black
- Recognize that primitive has non-zero width
- even thinnest line is 1 pixel thick
- Consider line as (thin) rectangle
- covers different (square) pixels to different extent
- In most cases should not set a single pixel to black
- Set intensity of pixel differently for each pixel covered
- Only horizontal and vertical lines effect only 1 pixel per row


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\section*{Properties of Unweighted Area Sampling}
- 1. Intensity decreases with increasing distance from pixel to edge
- 2. Primitives do not influence pixel they do not intersect
- 3. Equal areas contribute equal intensity
- distance from pixel center to area overlapped
- small area in corner contributes same as equal-sized area in center

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}

\section*{Polygon Aliasing}
- Aliasing problems can be serious for polygons
- Jaggedness of edges
- Small polygons neglected - Need compositing so color of one polygon does not totally determine color of pixel


All three polygons should contribute to color
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[^0]:    Angel: Interactive Computer Graphics 4E © Addison-Wesley 2005 KENT STATE ${ }^{4}$

[^1]:    Angel: Interactive Computer Graphics 4E © Addison-Wesley 2005
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[^2]:    Angel: Interactive Computer Graphics 4E © Addison-Wesley 2005

[^3]:    Angel: Interactive Computer Graphics 4E © Addison-Wesley 2005 KENT STATE ${ }^{24}$

[^4]:    Angel: Interactive Computer Graphics 4E © Addison-Wesley 2005

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