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Chapter 5: Threads

5.1

- Overview
- Multithreading Models
- Threading Issues
- Pthreads
- Solaris 2 Threads
- Windows 2000 Threads
- Linux Threads
- Java Threads

Operating System Concepts

Single and Multithreaded Processes Threads are lightweight processes – have own thread ID, PC, registers, stack code data files code data files registers registers registers registers stack stack stack stack thread single-threaded multithreaded **Operating System Concepts** Silberschatz, Galvin and Gagne ©2002 5.2



Benefits

- Responsiveness
 - Can run even if one thread blocked or busy
 - Web browser example one thread per client
- Resource Sharing
- Economy
 - Creating and context switching threads is low cost
 - Solaris 2: creating 30x, context switch 5x slower for procs

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- Utilization of MP Architectures
 - Run each thread on different CPU

Operating System Concepts



User Threads

- Thread management done by user-level threads library
- No need for kernel intervention
- Drawback : all may run in single process. If one blocks, all block.

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- Examples
 - POSIX Pthreads
 - Mach C-threads
 - Solaris threads

Operating System Concepts



Kernel Threads

- Supported by the Kernel
- Generally slower to create than user threads
- If one blocks another in the application can be run
- Can be scheduled on different CPUs in multiprocessor
- Examples
 - Windows 95/98/NT/2000
 - Solaris
 - Tru64 UNIX
 - BeOS
 - Linux

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Multithreading Models ■ Many-to-One ■ One-to-One ■ Many-to-Many Operating System Concepts 5.6 Silberschatz, Galvin and Gagne ©2002

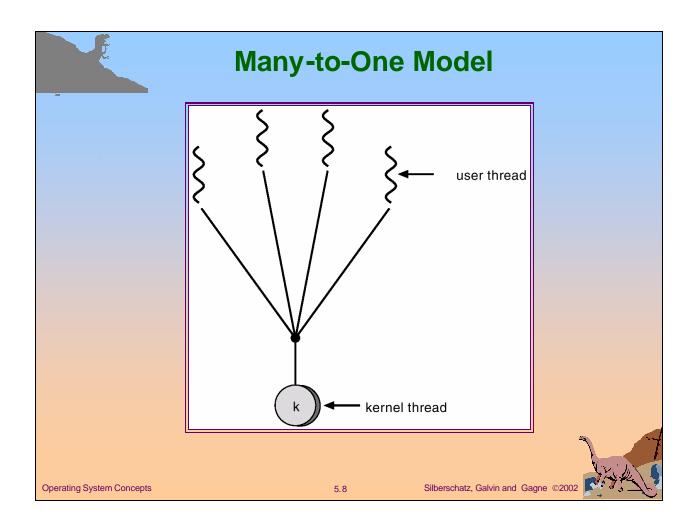


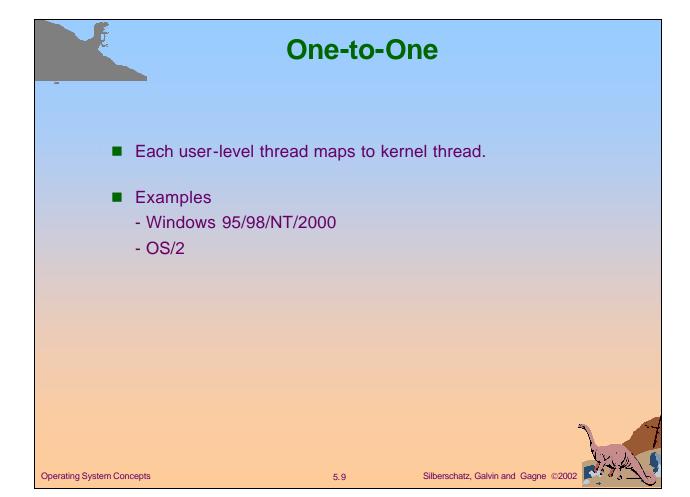
Many-to-One

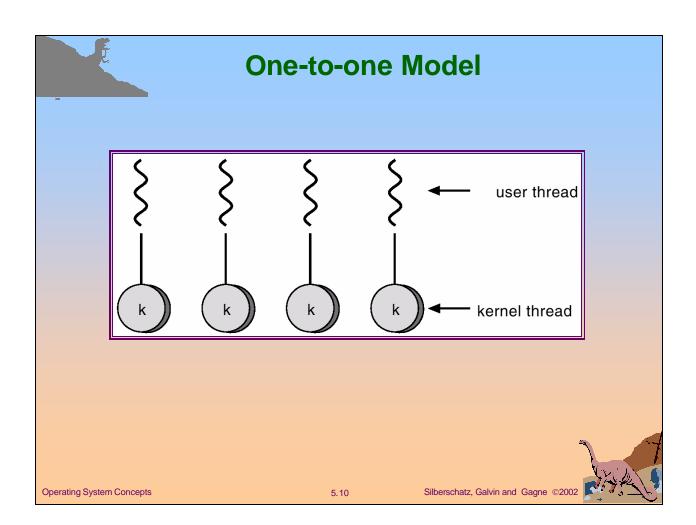
- Many user-level threads mapped to single kernel thread.
- Used on systems that do not support kernel threads.

Operating System Concepts

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Many-to-Many Model

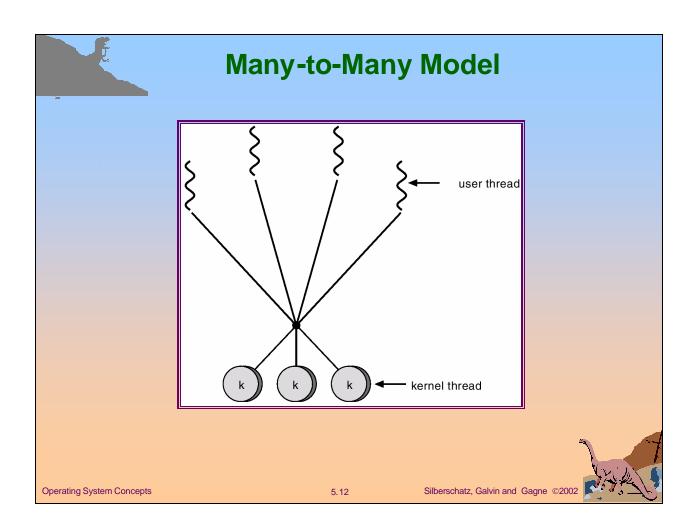
- Allows many user level threads to be mapped to many kernel threads.
- Allows the operating system to create a sufficient number of kernel threads.
- Solaris 2
- IRIX

Operating System Concepts

- HP-UX
- Tru64 Unix
- Windows NT/2000 with the *ThreadFiber* package

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Threading Issues

- Semantics of fork() and exec() system calls
 - Duplicate all threads or not
 - Exec replaces all threads
 - If call exec next no need to duplicate all threads.
- Thread cancellation.
 - Asynchronous or deferred (target thread checks periodically)
 - Resource reclamation problem
- Thread pools
 - Create pool of threads to do work
 - When server receives request awakens thread. Returns on finish.
 - Advantages:
 - Faster than creating threads
 - Limits number of threads in server and hence load on CPU

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■ Thread specific data

Operating System Concepts

Silberschatz,

Threading Issues

- Signal handling
 - Signals can be synchronous (e.g. illegal memory access) or asynchronous (e.g. i/o completion, ^C)
 - Handled by default handler or user-defined handler
 - Where should the thread be delivered?
 - To thread to which applies (synchronous signals)

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- To all threads in process
- To certain threads in process
- Assign a specific thread to receive all signals (Solaris 2)

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Pthreads

- a POSIX standard (IEEE 1003.1c) API for thread creation and synchronization.
- API specifies behavior of the thread library, implementation is up to developer of the library.
- Common in UNIX operating systems.

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Pthreads example

```
#include <pthread.h>
        #include <stdio.h>
                                     /* this data is shared by the thread(s) */
        int sum:
        void *runner(void *param);
                                                               /* the thread */
        main(int argc, char *argv[])
        pthread t tid;
                                 /* the thread identifier */
        pthread attr tattr; /* set of attributes for the thread */
        if (argc != 2) {
                  fprintf(stderr,"usage: a.out <integer value>\n");
                  exit();
        if (atoi(argv[1]) < 0) {
                  fprintf(stderr,"Argument %d must be non-negative\n",atoi(argv[1]));
                  exit();
        pthread_attr_init(&attr);
                                                    /* get the default attributes */
        pthread_create(&tid,&attr,runner,argv[1]);
                                                           /* create the thread */
                                             /* now wait for the thread to exit */
         pthread_join(tid,NULL);
        printf("sum = %d\n",sum);
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                                                            Silberschatz, Galvin and Gagne ©2002
                                             5.16
```

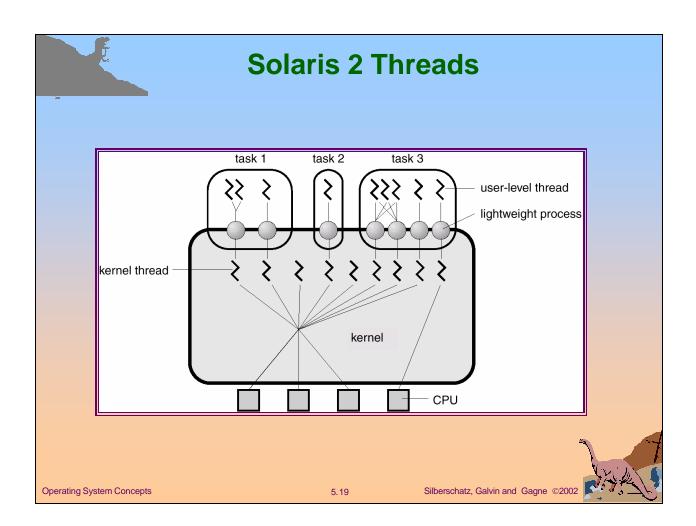


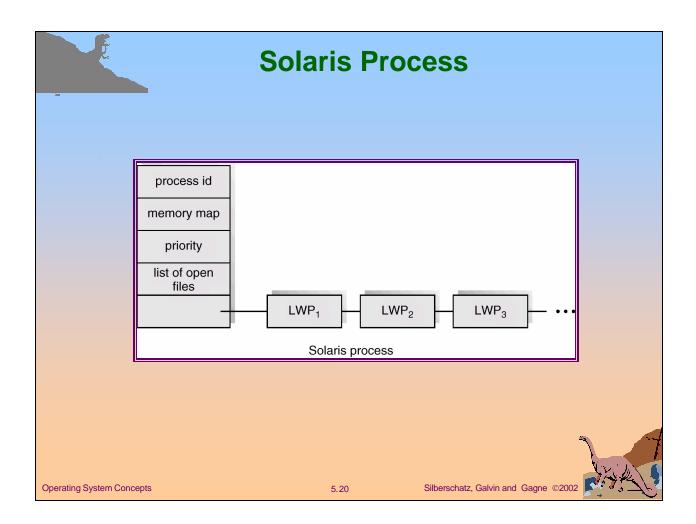

- User and Kernel level threads, Light weight processes (LWP)
- Process : one or more LWPs
- Each LWP has kernel thread
- One LWP is needed for each user thread that may block
- If kernel thread blocks, LWP, and user level thread also block
- If all LWPs in process block, but there are user level threads which could run, kernel creates new LWP
- Kernel "ages" LWPs and deletes unused ones after +-5 min
- Kernel threads may be bound to particular CPU



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Windows 2000 Threads

- Implements the one-to-one mapping.
- Each thread contains
 - a thread id
 - register set
 - separate user and kernel stacks
 - private data storage area

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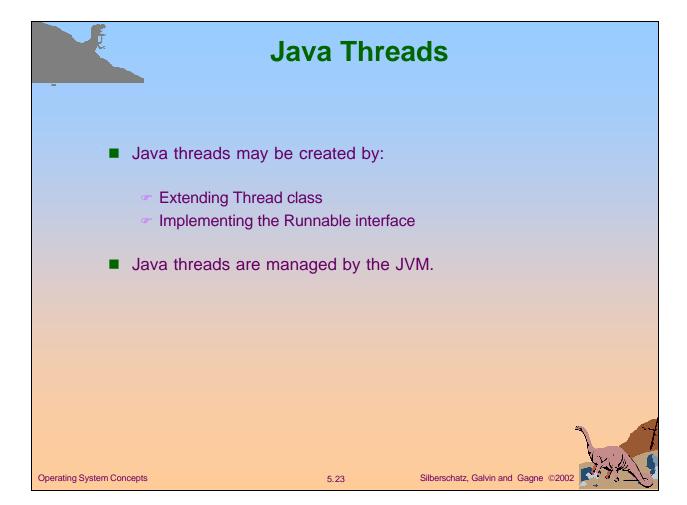


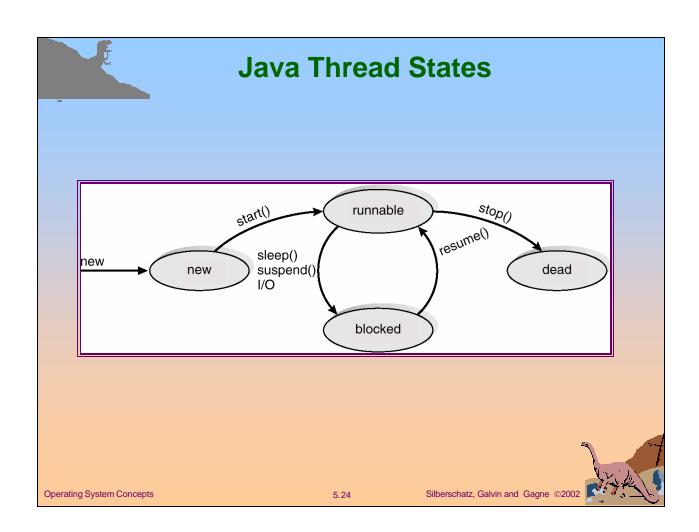
Linux Threads

- Linux refers to them as tasks rather than threads.
- Thread creation is done through clone() system call.
- Clone() allows a child task to share the address space of the parent task (process)
- The amount of parent process shared is determined by a set of flags passed as parameter in clone() call
 - None set, no sharing clone() is fork()
 - All set, everything shared

Operating System Concepts

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