

<b>CS 4/56101</b>	<b>Kent State University</b> Dept. of Math & Computer Science <u>LECT-11</u>
<b>Design and Analysis of Algorithms</b>	

# Splay Tree

# Splay Tree

- Idea:
  - in a Binary Search Tree or even in an AVL tree where the new records are going?
  - The later a key/record is added the higher will be its access time!
  - Consider a hospital situation..
- Motivation
  - How can we keep frequently accessed/ recently used keys near the root?



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# Splay Tree

- A Self-adjusting Tree
  - A Binary Search Tree
  - Not so well balanced explicitly.
- Adjustments
  - Every time a new node is inserted, it is positioned at the root.
  - Every time a node is retrieved, it becomes the root.
- Performance
  - A sequence of  $m$  insertions or retrievals with splaying a binary search tree of size  $n$  will never need more than  $m(1+3 \log n) + \log n$  upward moves of a target node.

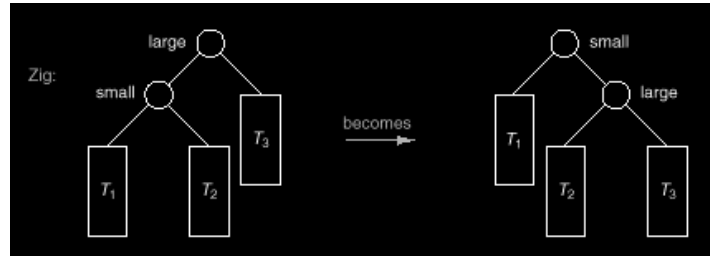


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## Idea of Zig and Zag

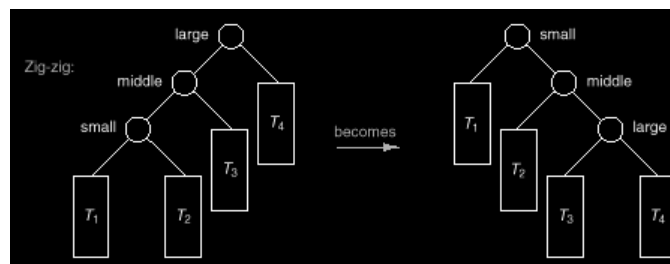
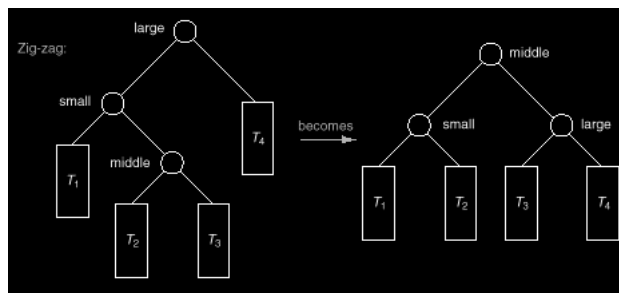
- Right-Rotation = Zig
- Left Rotation = Zag



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## Zig-Zag and Zig Zig



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## Basic Algorithm

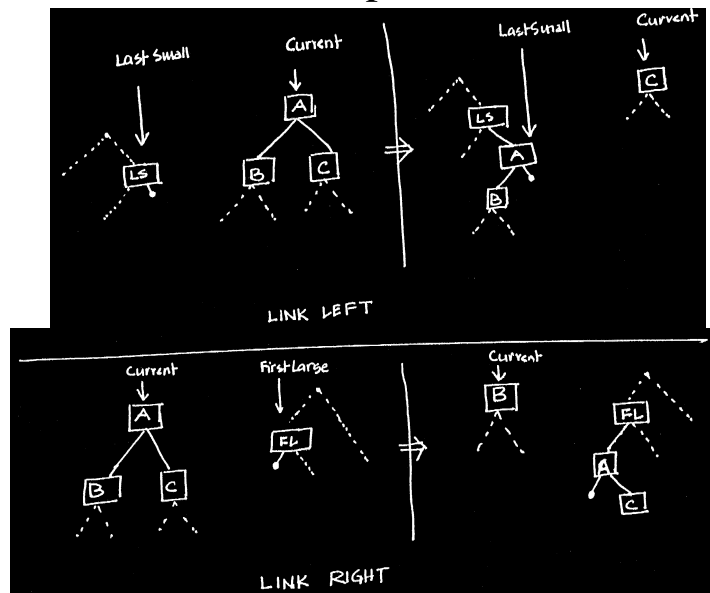
- Start from the root and keep on splitting the tree applying zig/zag and eventually bringing the target at the root.
- 
- Divide the tree in three parts:
  - Tree#1: all the nodes confirmed greater than the key.
  - Tree#3: all the nodes confirmed larger than the key.
  - Tree#2: if the key is there it must be inside this tree.



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## Some Operations



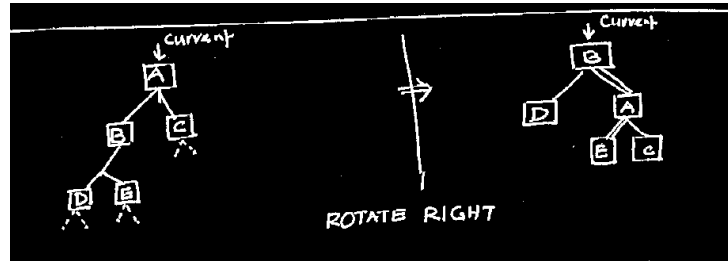
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## Some Operations (continued..)



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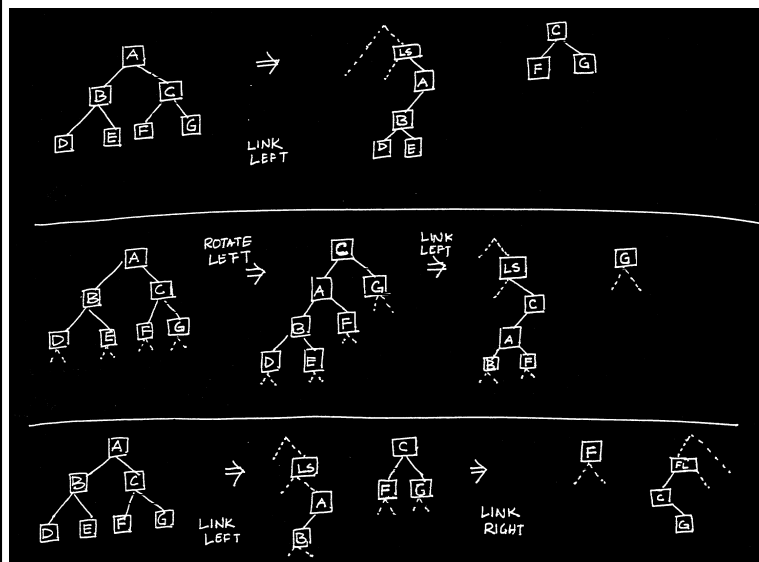


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## Algorithm



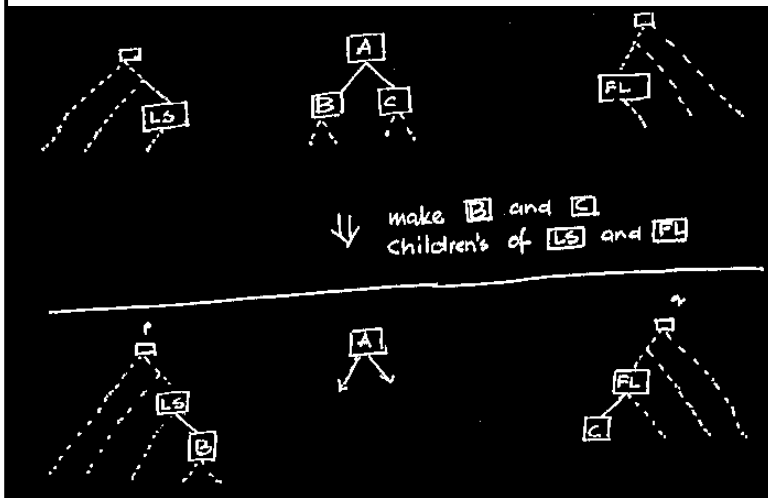
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Case  $I < C$

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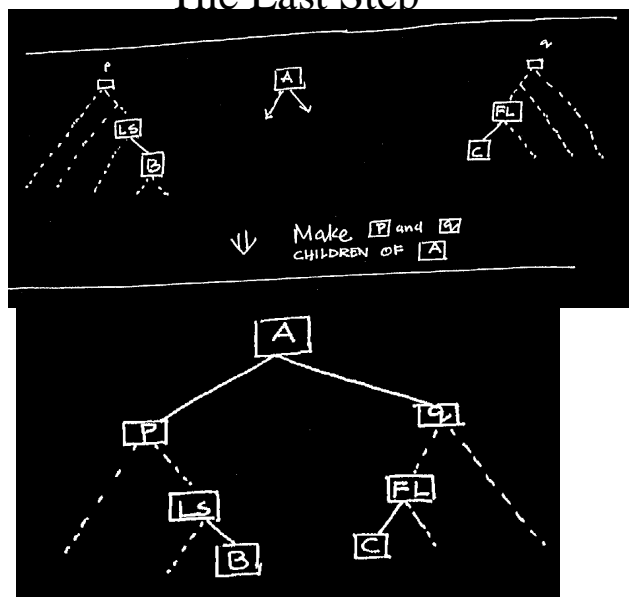
## Final Two Steps..



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## The Last Step

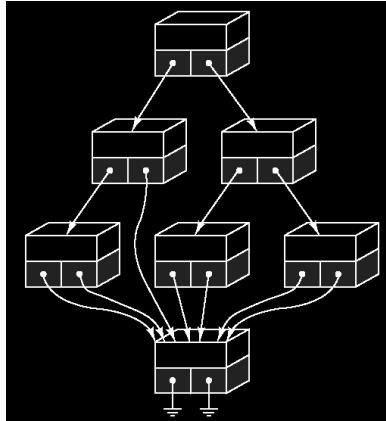


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## Sentinel Binary Tree

- No Need to Check for NULL Pointer. When we are at "Sentinel" node, we know we are at the leaf.
- In Splay Tree we start by copying the "target" at the sentinel node..
- The two pointers of the sentinel keeps track of the low and high sub-trees.



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```
TreeNode *TreeSplay(TreeNode *root, KeyType target)
{
    TreeNode *current;      /* the current position in the tree */
    TreeNode *child;        /* one of the children of current */
    TreeNode *lastsmall;    /* largest key known to be less than the target */
    TreeNode *firstlarge;   /* smallest key known to be greater than the target */
    extern TreeNode *sentinel;

```

```
    sentinel->entry.key = target; /* Establish sentinel for searching. */
    lastsmall = firstlarge = sentinel;

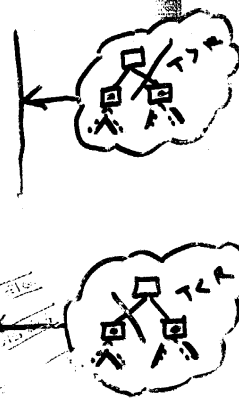
```

```
    for (current = root; NE(current->entry.key, target); )
    {
        if (LT(current->entry.key, target)) {
            child = current->right;
            if (EQ(target, child->entry.key)) {
                current = LinkLeft(current, &lastsmall);
            } else if (GT(target, child->entry.key)) {
                current = RotateLeft(current);
                current = LinkLeft(current, &lastsmall);
            } else {
                current = LinkLeft(current, &lastsmall);
                current = LinkRight(current, &firstlarge);
            }
        } else {
            child = current->left;
            if (EQ(target, child->entry.key)) {
                current = LinkRight(current, &firstlarge);
            } else if (LT(target, child->entry.key)) {
                current = RotateRight(current);
                current = LinkRight(current, &firstlarge);
            } else {
                current = LinkRight(current, &firstlarge);
                current = LinkLeft(current, &lastsmall);
            }
        }
    }

```

```
    if (current == sentinel) {
        printf("Target has been inserted as root of the tree.");
        root = current = MakeNode(target, sentinel);
    }

```



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