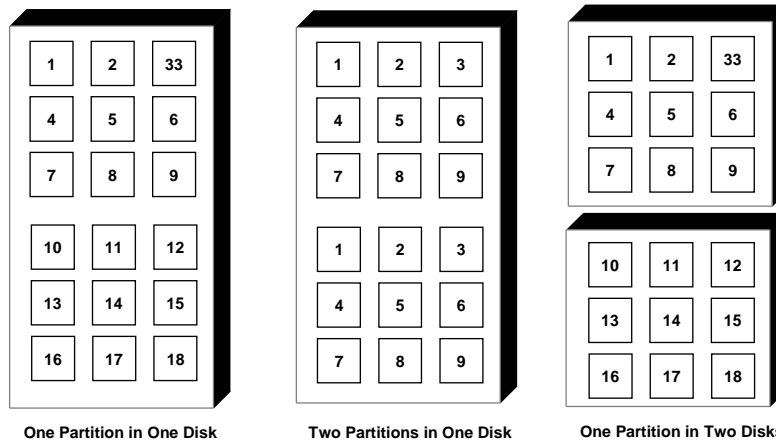


## File System Concepts

- **What is a File?**
- **Attributes of A File**
- **Operations on a File**
- **Access Modes**
  - ◆ Sequential vs. direct
- **Directory**
  - ◆ Tree vs. acyclic-graph
- **Protection**
- **Shared Files**
- **File Mounting**
- **Open File Table**

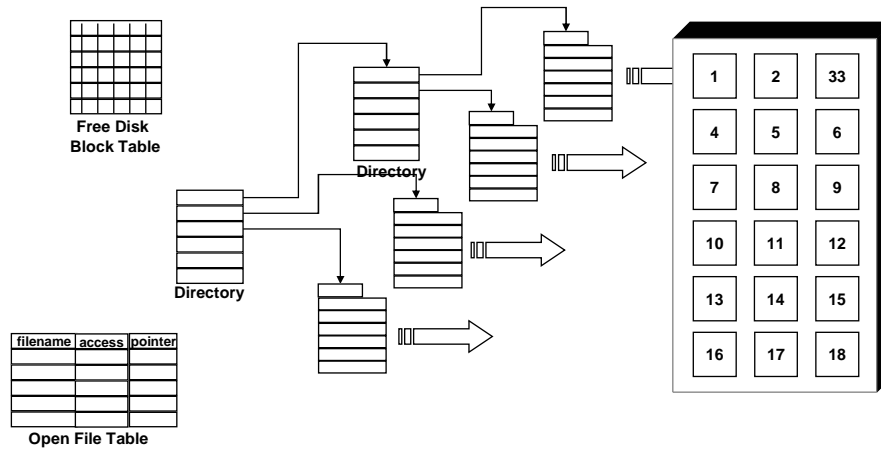
Os-slide#1

## File System Organization: Partition and Physical Disks



Os-slide#2

## Organization of a Partition



Os-slide#3

## Contiguous Allocation

A Contiguous set of block are allocated for each file. First-Fit Best-Fit can be used.

### Advantages:

- No space needed for inodes.
- Fast sequential and direct access.

### Disadvantages:

- External Fragmentation.
- No dynamic growth. File size must be known ahead.

Os-slide#4

## Linked Allocation

Blocks of a file are scattered in disk. But each block has a pointer to next block at the end.

Advantages:

- No external fragmentation.
- Dynamically blocks can be added and deleted.

Disadvantages:

- 4 bytes in each block is needed. ( $4/512 = 12\%$  overhead)
- Direct access can be very slow! Reading  $n$ th block need reading  $n$  blocks.
- Reliability! Consider losing one block in a chain!

Os-slide#5

## FAT based Linked Allocation

The first block of a partition contains a table (File Allocation Table) implementation of linked list. Used by MS-DOS

Advantages:

- Reliability Improved.
- Direct access time is better.

Os-slide#6

## Indexed Allocation

Each file has its own index table block.

Advantages:

- No external fragmentation.
- Dynamically blocks can be added and deleted.
- Fast direct and sequential access.

Disadvantages:

- Space wasted for index blocks for small files.
- File size is limited by block size.

BSD Unix approach:

- Combined scheme.

Os-slide#7

## Free Space Management

• Bit Vectors:

- ◆ Very efficient, but may be very large for large disks (1.3 GB disk with .5K blocks need 310Kbytes).

• Linked List:

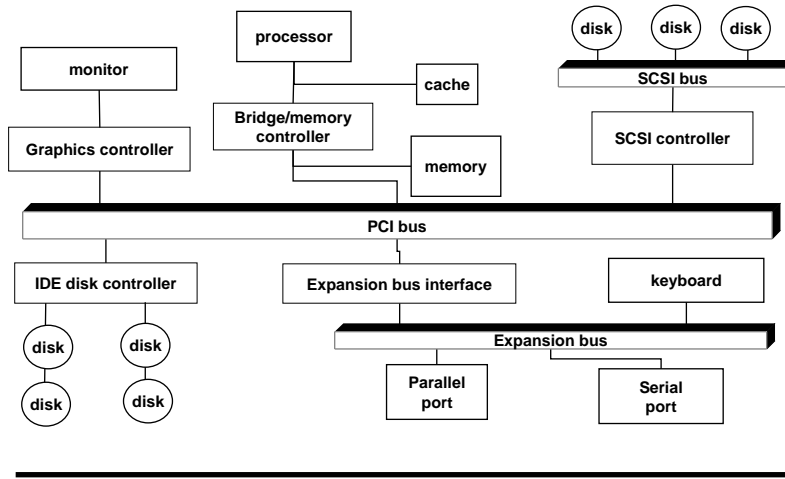
- ◆ Each free block contains pointer to the next one. Takes a disk access to find a free block.

• Grouping:

- ◆ The index of first n free blocks are in the first block. A large number of free blocks can be found at once.

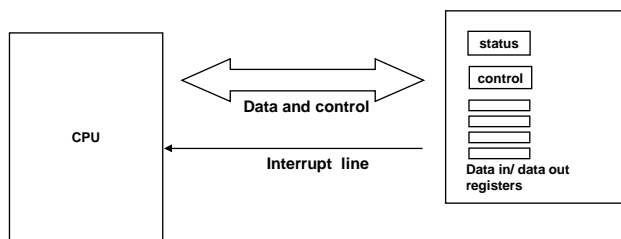
Os-slide#8

## A Typical PC bus Architecture



Os-slide#9

## Communication between CPU and Controller



**I/O vs. Memory Mapped**  
**Polling**  
**Interrupt Vector**  
**DMA**  
**Spooling**

- Character stream vs. block devices
- Sequential vs. random access
- Synchronous vs. asynchronous devices
- Sharable vs. dedicated devices
- Slow vs. fast devices
- Read-only, write-only, read/write devices

Os-slide#10

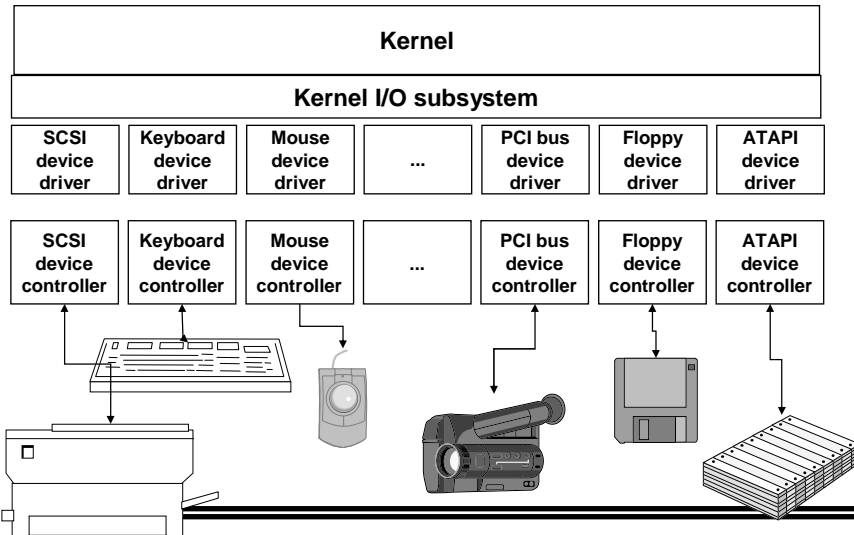
# Disk Scheduling

Init Position 53

Request 98, 183, 37, 122, 14, 124, 65, 67

- **FCFS Algorithm:**
  - ◆ Read Order: (53) 98, 183, 37, 122, 14, 124, 65, 67
- **SSTF Algorithm:**
  - ◆ Read Order: (53) 65, 67, 98, 122, 124, 183
- **SCAN Algorithm:**
  - ◆ Read Order: (53) 37, 14, (0), 65, 67, 98, 122, 124, 183
- **C-SCAN Algorithm:**
  - ◆ Read Order: (53) 37, 14, (0), 183, 124, 122, 98, 67, 65
- **LOOK Algorithm:**
- **C-LOOK Algorithm:**

# OS Organization



## **Disk Management**

- **Formatting**
- **Boot Block**
- **Bad Blocks**
- **Swap Space Location**
- **Swap Space Management**