

Representing / Encoding Integers

- Encoding = symbolic representation of a value, in some specified number of digits, in some specified alphabet (here, 0 & 1)
- *Unsigned binary* representation
 - Represent only positive numbers
 - Direct binary representation
 - Examples (8-bit unsigned binary)

$$+13_{10} = +1101_2 = 00001101_{2ub}$$
 - Can represent 0 to 255 in 8 bits
- *Signed magnitude* representation
 - Precede number with sign bit
 - 0 = positive, 1 = negative
 - Examples (8-bit signed magnitude)

$$+13_{10} = +1101_2 = 00001101_{2sm}$$

$$-13_{10} = -1101_2 = 10001101_{2sm}$$
 - Can represent -127 to +127 in 8 bits

1

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Encoding Integers (cont.)

- *Excess n* representation
 - Add *bias value* (n) to number, then represent directly in binary
 - Examples (8-bit excess 128)

$$+13_{10} = 13 + 128 = 141 = 10001101_{2ex128}$$

$$-13_{10} = -13 + 128 = 115 = 01110011_{2ex128}$$
 - Can represent -128 to 127 in 8 bits
- *Two's complement* representation
 - Represent positive numbers in n -bit signed magnitude form
 - Represent negative numbers as $2^n - N$
 - Examples (8-bit two's complement)

$$+13_{10} = +0001101_2 = 00001101_{2'scomp}$$

$$-13_{10} = 256 - 13 = 243 = 11110011_{2'scomp}$$
 - Can represent -128 to 127 in 8 bits

2

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8-Bit Numerical Representations

| Numeric Value | Unsigned Binary | Signed Magnitude | Excess 128 | Excess 127 | Two's Complement |
|---------------|-----------------|------------------|------------|------------|------------------|
| 255 | 11111111 | N/A | N/A | N/A | N/A |
| 254 | 11111110 | | | | |
| ... | ... | | | | |
| 128 | 10000000 | | | 11111111 | |
| 127 | 01111111 | 01111111 | 11111111 | 11111110 | 01111111 |
| 126 | 01111110 | 01111110 | 11111110 | 11111101 | 01111110 |
| ... | ... | ... | ... | ... | ... |
| 2 | 00000010 | 00000010 | 10000010 | 10000001 | 00000010 |
| 1 | 00000001 | 00000001 | 10000001 | 10000000 | 00000001 |
| 0 | 00000000 | 00000000 | 10000000 | 01111111 | 00000000 |
| -1 | N/A | 10000001 | 01111111 | 01111110 | 11111111 |
| -2 | | 10000010 | 01111110 | 01111101 | 11111110 |
| ... | | ... | ... | ... | ... |
| -127 | | 11111111 | 00000001 | 00000000 | 10000001 |
| -128 | | N/A | 00000000 | N/A | 10000000 |

3

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Worksheet — Two's Complement Addition and Subtraction

- Perform the following arithmetic operations in 4-bit two's complement arithmetic, showing your work:

$$2 + 3 =$$

$$2 - 4 =$$

$$(-7) - (-5) =$$

$$(-5) + (-2) =$$
- To add two numbers, add as usual
- To subtract two numbers, negate the second, and then add
 - Two ways to negate a two's complement number:
 - Invert all the bits, and add 1 to the result
 - Scan right to left, keep all bit the same, but invert all bits after passing the first "1"
 - In 8 bits: $-13 = -00001101 = 11110011$
 - Why subtract this way??

4

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Fixed-Point Representation

- Suppose we have a number of values to represent. One way to do so would be use a specific number of digits, assuming a decimal point to the right of the lsd:

000500. 006000. 010000.

- Or we could assume that the decimal point is further to the right of the lsd:

000005 __. 000060 __. 000100 __.

- We could also assume that the decimal point is between the msd and lsd:

00500.0 06000.0 10000.0

- All of these are called *fixed-point* representations — ones where the decimal point is fixed at one specific place for encoding all numbers

5

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Floating-Point Representation

- Consider the following ways of representing the decimal values:

| | | | |
|------|--------|---------|--------|
| 5 | 60 | 100 | (x100) |
| 500 | 6,000 | 10,000 | (x1) |
| 5000 | 60,000 | 100,000 | (x0.1) |

- Internally, we can represent a value as:

number, n multiplied by *scale factor, s*

| | | | |
|--------|--------|--------|-----------|
| 000005 | 000060 | 000100 | $s = 100$ |
| 000500 | 006000 | 010000 | $s = 1$ |
| 005000 | 060000 | 100000 | $s = 0.1$ |

- Encoding a value using a number (*mantissa*) and an scale factor (*characteristic*, or *exponent*, to an implicit base) is called a *floating-point* representation

6

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Floating-Point Representation

- In general, this is all very similar to *scientific notation* — writing a number as as the product of a decimal number and a power of 10

- 5×10^2 60×10^2 100×10^2
- 500×10^0 6000×10^0 100000×10^0

- A number can be represented uniquely in *normal form* — where we insist that there be exactly one non-zero digit to the left of the decimal point

- 5×10^2 6×10^3 1×10^4

- All significant digits must be included. Thus the following three values are different:

- 5×10^2 5.0×10^2 5.000×10^2

7

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Homework #5 — Due 11/9/98 (Part 1)

1. Perform the following arithmetic operations in 5-bit two's complement arithmetic, showing your work:

$$6 + 7 =$$

$$6 - 7 =$$

$$6 + (-4) =$$

$$(-2) - (-5) =$$

8

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