What is an Operating System? (Review)

- An operating system (OS) is the interface between the user and the hardware
 - It implements a virtual machine that is easier to program than bare hardware
- An OS provides standard services (an interface) which are implemented on the hardware, including:
 - Processes, CPU scheduling, memory management, file system, networking
- The OS coordinates multiple applications and users (multiple processes) in a fair and efficient manner
- The goal in OS development is to make the machine convenient to use (a software engineering problem) and efficient (a system and engineering problem)

History of Operating Systems (cont.)

- Phase 1 hardware is expensive, humans are cheap
 - 3. Overlapped CPU & I/O operations
 - First: buffer slow I/O onto fast tape drives connected to CPU, replicate I/O devices
 - Later: *spool* data to disk
 - 4. Multiprogrammed batch systems
 - Multiple jobs are on the disk, waiting to run
 - Multiprogramming run <u>several</u> programs at the "same" time
 - Pick some jobs to run (*scheduling*), and put them in memory (*memory management*)
 - Run one job; when it waits on something (tape to be mounted, key to be pressed), switch to another job in memory
 - First big failures:

3

- MULTICS announced in 1963, not released until 1969
- IBM's OS/360 released with 1000 known bugs
- OS design should be a science, not an art

Fall 2000, Lecture 02

2

History of Operating Systems

- Phase 0 hardware is a very expensive experiment; no operating systems exist
 - 1. One user at console
 - One function at a time (computation, I/O, user think/response)
 - Program loaded via card deck
 Libraries of device drivers (for I/O)
 - User debugs at console
- Phase 1 hardware is expensive, humans are cheap
 - 2. Simple batch processing: load program, run, print results, dump, repeat
 - User gives program (cards or tape) to the operator, who schedules the jobs
 - Resident monitor automatically loads, runs, dumps user jobs
 - Requires memory management (relocation) and protection
 - More efficient use of hardware, but debugging is more difficult (from dumps) Fall 2000, Lecture 02

History of Operating Systems (cont.)

- Phase 2 hardware is less expensive than before, humans are expensive
 - 5. Interactive timesharing
 - Lots of cheap terminals, one computer
 - All users interact with system at once
 - Debugging is much easier
 - Disks are cheap, so put programs and data online
 - 1 punch card = 100 bytes
 - 1MB = 10K cards
 - OS/360 was several feet of cards
 - New problems:
 - Need preemptive scheduling to maintain adequate response time
 - Need to avoid *thrashing* (swapping programs in and out of memory too often)
 - Need to provide adequate security measures
 - Success: UNIX developed at Bell Labs so a couple of computer nerds (Thompson, Ritchie) could play Star Trek on an unused PDP-7 minicomputer

History of Operating Systems (cont.)	History Lessons
 Phase 3 — hardware is very cheap, humans are expensive 6. Personal computing 	None of these operating systems were particularly bad; each depended on tradeoffs made at that point in time
CPUs are cheap enough to put one in	 Technology changes drive OS changes
 each terminal, yet powerful enough to be useful Computers for the masses! Return to simplicity; make OS simpler by getting rid of support for multiprogramming, concurrency, and protection 	 Since 1953, there has been about 9 orders of magnitude of change in almost every computer system component Unprecedented! In past 200 years, gone from horseback (10 mph) to Concorde
	(1000 mph), only 2 orders of magnitude
Modern operating systems are:	Changes in "typical" academic computer:
 Enormous Small OS = 100K lines of code Big OS = 10M lines Complex (100-1000 person year of work) 	1981 1996 MIPS 1 400 price / MIPS \$100,000 \$50 memory 128 KByte 64 MByte disk 10 MByte 4 GByte
 Poorly understood (outlives its creators, too large for one person to comprehend) Fall 2000, Lecture 02 	network 9600 bit/sec 155 Mb/s address bits 16 64
Modern OS Functionality (Review)	More Recent Developments
•	More Recent Developments Parallel operating systems
(Review)	
 (Review) Concurrency Multiple processes active at once Processes can communicate 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled
 (Review) Concurrency Multiple processes active at once 	 Parallel operating systems Shared memory, shared clock
 (Review) Concurrency Multiple processes active at once Processes can communicate Processes may require mutually- 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled processors
 (Review) Concurrency Multiple processes active at once Processes can communicate Processes may require mutually- exclusive access to some resource 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled processors Appearance of single operating system
 (Review) Concurrency Multiple processes active at once Processes can communicate Processes may require mutually- exclusive access to some resource CPU scheduling, resource management 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled processors Appearance of single operating system Distributed operating systems
 (Review) Concurrency Multiple processes active at once Processes can communicate Processes may require mutually-exclusive access to some resource CPU scheduling, resource management Memory management — allocate memory to processes, move processes 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled processors Appearance of single operating system Distributed operating systems No shared memory, no shared clock Small number of loosely-coupled
 (Review) Concurrency Multiple processes active at once Processes can communicate Processes may require mutually-exclusive access to some resource CPU scheduling, resource management Memory management — allocate memory to processes, move processes between disk and memory File system — allocate space for storage of programs and data on disk Networks and distributed computing — 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled processors Appearance of single operating system Distributed operating systems No shared memory, no shared clock Small number of loosely-coupled processors Appearance of single operating system is ideal goal, but not realized in practice May try to simulate a shared memory
 (Review) Concurrency Multiple processes active at once Processes can communicate Processes may require mutually- exclusive access to some resource CPU scheduling, resource management Memory management — allocate memory to processes, move processes between disk and memory File system — allocate space for storage of programs and data on disk 	 Parallel operating systems Shared memory, shared clock Large number of tightly-coupled processors Appearance of single operating system Distributed operating systems No shared memory, no shared clock Small number of loosely-coupled processors Appearance of single operating system is ideal goal, but not realized in practice

8

7

Fall 2000, Lecture 02