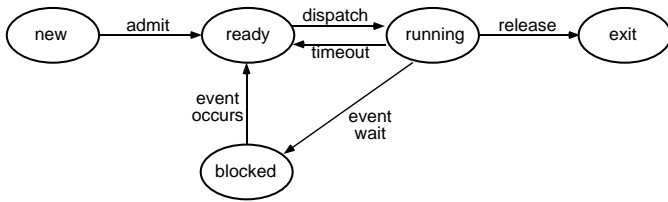


CPU Scheduling



- The *CPU scheduler* (sometimes called the *dispatcher* or *short-term scheduler*):
 - Selects a process from the ready queue and lets it run on the CPU
 - Assumes all processes are in memory, and one of those is executing on the CPU
 - Crucial in multiprogramming environment
 - Goal is to maximize CPU utilization
- *Non-preemptive* scheduling — scheduler executes only when:
 - Process is terminated
 - Process switches from running to blocked

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Process Execution Behavior

- Assumptions:
 - One process per user
 - One thread per process
 - Processes are independent, and compete for resources (including the CPU)
- Processes run in CPU - I/O burst cycle:
 - Compute for a while (on CPU)
 - Do some I/O
 - Continue these two repeatedly
- Two types of processes:
 - CPU-bound — does mostly computation (long CPU burst), and very little I/O
 - I/O-bound — does mostly I/O, and very little computation (short CPU burst)

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First-Come-First-Served (FCFS)

- Other names:
 - First-In-First-Out (FIFO)
 - Run-Until-Done
- Policy:
 - Choose process from ready queue in the order of its arrival, and run that process non-preemptively
 - Early FCFS schedulers were overly non-preemptive: the process did not relinquish the CPU until it was finished, even when it was doing I/O
 - Now, non-preemptive means the scheduler chooses another process when the first one terminates or blocks
- Implement using FIFO queue (add to tail, take from head)
- Used in Nachos (as distributed)

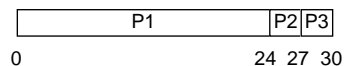
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FCFS Example

- Example 1:

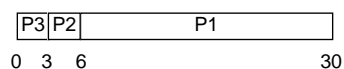
Process (Arrival Order)	P1	P2	P3
Burst Time	24	3	3
Arrival Time	0	0	0



$$\text{average waiting time} = (0 + 24 + 27) / 3 = 17$$

- Example 2:

Process (Arrival Order)	P3	P2	P1
Burst Time	3	3	24
Arrival Time	0	0	0



$$\text{average waiting time} = (0 + 3 + 6) / 3 = 3$$

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Scheduling in Nachos (Review)

- `main()` (in `threads/main.cc`)
calls `Initialize()` (in `threads/system.cc`)
 - which starts scheduler, an instance of class `Scheduler` (defined in `threads/scheduler.h`, `scheduler.cc`)
- Interesting functions:
 - Mechanics of running a thread:
 - `Scheduler::ReadyToRun()` — puts a thread at the tail of the ready queue
 - `Scheduler::FindNextToRun()` — returns thread at the head of the ready queue
 - `Scheduler::Run()` — switches to thread
 - Scheduler is non-preemptive FCFS — chooses next process when:
 - Current thread terminates
 - Current thread calls `Thread::Yield()` to explicitly yield the CPU
 - Current thread calls `Thread::Sleep()` (to block (wait) on some event)

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Scheduling in Nachos (Review)

```
Scheduler::Scheduler ( )
{
    readyList = new List;
}

void
Scheduler::ReadyToRun (Thread *thread)
{
    DEBUG('t',
        "Putting thread %s on ready list.\n",
        thread->getName());
    thread->setStatus(READY);
    readyList->Append((void *)thread);
}

Thread *
Scheduler::FindNextToRun ( )
{
    return (Thread *)readyList->Remove();
}
```

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Scheduling in Nachos (Review)

```
void
Scheduler::Run (Thread *nextThread)
{
    Thread *oldThread = currentThread;

    oldThread->CheckOverflow();
    currentThread = nextThread;
    currentThread->setStatus(RUNNING);

    DEBUG('t', "Switching from thread \"%s\"
to thread \"%s\"\n",oldThread->getName(),
    nextThread->getName());
    SWITCH(oldThread, nextThread);
    DEBUG('t', "Now in thread \"%s\"\n",
    currentThread->getName());

    if (threadToBeDestroyed != NULL) {
        delete threadToBeDestroyed;
        threadToBeDestroyed = NULL;
    }
}
```

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Manipulating Threads in Nachos (Review)

```
void
Thread::Fork(VoidFunctionPtr func, int arg)
{
    DEBUG('t',"Forking thread \"%s\" with
    func = 0x%x, arg = %d\n",
    name, (int) func, arg);

    StackAllocate(func, arg);

    IntStatus oldLevel = interrupt->
    SetLevel(IntOff);
    scheduler->ReadyToRun(this);
    (void) interrupt->SetLevel(oldLevel);
}
```

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Manipulating Threads in Nachos (cont.)

```
void
Thread::Yield ()
{
    Thread *nextThread;

    IntStatus oldLevel = interrupt->
        SetLevel(IntOff);

    ASSERT(this == currentThread);
    DEBUG('t', "Yielding thread \"%s\"\n",
        getName());

    nextThread = scheduler->
        FindNextToRun();
    if (nextThread != NULL) {
        scheduler->ReadyToRun(this);
        scheduler->Run(nextThread);
    }
    (void) interrupt->SetLevel(oldLevel);
}
```

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Manipulating Threads in Nachos (cont.)

```
void
Thread::Sleep ()
{
    Thread *nextThread;

    ASSERT(this == currentThread);
    ASSERT(interrupt->getLevel() == IntOff);
    DEBUG('t', "Sleeping thread \"%s\"\n",
        getName());

    status = BLOCKED;
    while ((nextThread = scheduler->
        FindNextToRun()) == NULL)
        interrupt->Idle();

    scheduler->Run(nextThread);
}
```

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Semaphores in Nachos (Review)

```
void
Semaphore::P()
{
    IntStatus oldLevel = interrupt->
        SetLevel(IntOff); // disable interrupts

    while (value == 0) { // sema not avail
        queue-> // so go to sleep
            Append((void *)currentThread);
        currentThread->Sleep();
    }

    value--; // semaphore available,
            // consume its value

    (void) interrupt-> // re-enable interrupts
        SetLevel(oldLevel);
}
```

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Semaphores in Nachos (cont.) (Review)

```
void
Semaphore::V()
{
    Thread *thread;

    IntStatus oldLevel = interrupt->
        SetLevel(IntOff);

    thread = (Thread *)queue->Remove();
    if (thread != NULL) // make thread ready,
        // consuming the V immediately
        scheduler->ReadyToRun(thread);

    value++;

    (void) interrupt->SetLevel(oldLevel);
}
```

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