Deadlock

- Consider this example:

Process A	Process B
printer–>wait( );	disk–>wait( );
disk->wait( );	printer->wait( );
print file	*print file*
printer->signal( );	disk->signal( );
disk->signal( );	printer->signal( );

- *Deadlock* occurs when two or more processes are each waiting for an event that will never occur, since it can only be generated by another process in that set
- Deadlock is one of the more difficult problems that OS designers face
  - As we examine various approaches to dealing with deadlock, notice the tradeoffs between how well the approach solves the problem, and its performance /OS overhead

Deadlock (cont.)

- OS must distribute system *resources* among competing processes:
  - CPU cycles — preemptable
  - Memory space — preemptable
  - Files — non-preemptable
  - I/O devices (printer) — non-preemptable
- A request for a type of resource can be satisfied by any resource of that type
  - Use any 100 bytes in memory
  - Use either one of two identical printers
- Process *requests* resource(s), *uses* it/them, then *releases* it/them
  - We will assume here that the resource is *re-usable*; it is not *consumed*
  - Waits if resource is not currently available

Deadlock Conditions

- These 4 conditions are **necessary** and **sufficient** for deadlock to occur:
  - **Mutual exclusion** — if one process holds a resource, other processes requesting that resource must wait until the process releases it (only one can use it at a time)
  - **No preemption** — resources are released voluntarily; neither another process nor the OS can force a process to release a resource
  - **Hold and wait** — processes are allowed to *hold* one (or more) resource and be *waiting* to acquire additional resources that are being held by other processes
  - **Circular wait** — there must exist a set of waiting processes such that P0 is waiting for a resource held by P1, P1 is waiting for a resource held by P2, … Pn-1 is waiting for a resource held by Pn, and Pn is waiting for a resource held P0

Deadlock Prevention

- Basic idea: ensure that one of the 4 conditions for deadlock can not hold
- **Mutual exclusion** — if one process holds a resource, other processes requesting that resource must wait until the process releases it
  - Hard to avoid mutual exclusion for non-sharable resources
    - Printers
    - Files
    - I/O devices or network connections
  - For printer, avoid mutual exclusion through spooling — then process won't have to wait on physical printer
  - However, many resources are sharable, so deadlock can be avoided for them
    - Read-only files (binaries, perhaps)
    - Most files in your account

Deadlock Prevention (cont.)

- **No preemption** — resources are released voluntarily; neither another process nor the OS can force a process to release a resource
  - To avoid, allow preemption
    - If process A requests resources that aren't available, see who holds those resources
      - If the holder (process B) is waiting on additional resources, preempt the resource requested by process A
      - Otherwise, process A has to wait
        - » While waiting, some of its current resources may be preempted
        - » Can only wake up when it acquires the new resources plus any preempted resources
    - If a process requests a resource that can not be allocated to it, **all** resources held by that process are preempted
      - Can only wake up when it can acquire all the requested resources
    - Only works for resources whose state can be saved/restored (memory, not printer)

Deadlock Prevention (cont.)

- **Hold and wait** — processes are allowed to *hold* one (or more) resource and be *waiting* to acquire additional resources that are being held by other processes
  - To avoid, ensure that whenever a process requests a resource, it doesn't hold any other resources
    - Request all resources (at once) at beginning of process execution
      - Process which loops forever?
    - Request all resources (at once) at any point in the program
    - To get a new resource, release all current resources, then try to acquire new one plus old ones all at once
  - Difficult to know what to request in advance
  - Wasteful; ties up resources and reduces resource utilization
  - Starvation is possible

Deadlock Prevention (cont.)

- **Circular wait** — there must exist a set of waiting processes such that P0 is waiting for a resource held by P1, P1 is waiting for a resource held by P2, … Pn-1 is waiting for a resource held by Pn, and Pn is waiting for a resource held P0
  - To avoid, impose a total order on all resources, and require process to request resource in that order
    - Order: disk drive, printer, CDROM
    - Process A requests disk drive, then printer
    - Process B requests disk drive, then printer
    - Process B does <u>not</u> request printer, then disk drive, which could lead to deadlock
  - Order should be in the logical sequence that the resources are usually acquired
    - Allow process to release all resources, and start request sequence over
    - Or force process to request total number of each resource in a single request

Dealing with Deadlock

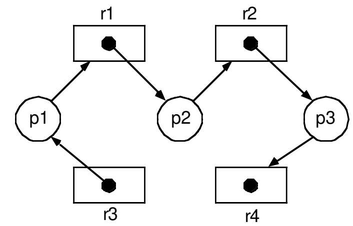
- *The Ostrich Approach* — stick your head in the sand and ignore the problem
- *Deadlock prevention* — prevent deadlock from occurring by eliminating one of the 4 deadlock conditions
- *Deadlock detection* algorithms — detect when deadlock has occurred
  - *Deadlock recovery* algorithms — break the deadlock
- *Deadlock avoidance* algorithms — consider resources currently available, resources allocated to each process, and possible future requests, and only fulfill requests that will not lead to deadlock

Resource-Allocation Graph

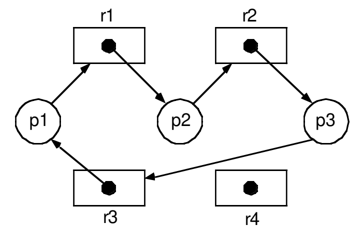
- The deadlock conditions can be modeled using a directed graph called a *resource-allocation graph* (RAG)
  - 2 kinds of nodes:
    - *Boxes* — represent resources
      - Instances of the resource are represented as dots within the box
    - *Circles* — represent processes
  - 2 kinds of (directed) edges:
    - *Request edge* — from process to resource — indicates the process has requested the resource, and is waiting to acquire it
    - *Assignment edge* — from resource instance to process — indicates the process is holding the resource instance
  - When a request is made, a request edge is added
    - When request is fulfilled, the request edge is transformed into an assignment edge
    - When process releases the resource, the assignment edge is deleted

Interpreting a RAG With Single Resource Instances

- If the graph does **not** contain a cycle, then **no** deadlock exists



- If the graph **does** contain a cycle, then a deadlock **does** exist



- With **single** resource instances, a **cycle** is a **necessary** and **sufficient** condition for deadlock