## **Supercomputing**

- Supercomputers have higher than average speed and capacity
- US government involvement
  - High Performance Computing and Communications Program (HPCC) federal agencies, industry, academia
  - DOE Accelerated Strategic Computing Initiative (ASCI)
    - Los Alamos National Labs & clusters of SMPs
    - One program three labs (LANL, ?, ?)
- Grand Challenge problems (1993)
  - Magnetic recording technology, high speed civil transports, catalysts, ocean modeling, digital anatomy, air pollution, Venus imaging

Fall 2000, Lecture 02

## Example Problem: Weather Prediction (in "Parallel Programming", by Wilkinson)

Atmosphere divided into 3D "cells", computations involve temperature, pressure, humidity, wind speed and direction, etc., are computed at time intervals, using info from previous interval

### ■ Some numbers:

- Suppose cell = 1 mile<sup>3</sup>, atmosphere modeled to height of 10 miles, gives 5 x 10<sup>8</sup> cells
- Suppose each calculation = 200 floating point operations, in one time step need 10<sup>11</sup> fp ops
- Suppose forecast weather over 10 days, at 10-minute intervals, would be 10<sup>4</sup> time steps and 10<sup>15</sup> fp ops
- Suppose computer runs at 100 Mflops (10<sup>8</sup> floating point operations per second), calcs would take 10<sup>7</sup> seconds = 100 days

Fall 2000, Lecture 02

Fall 2000, Lecture 02

# Two Taxonomies for Classifying Computer Systems

- Michael Flynn (1966)
  - SISD single instruction, single data
  - SIMD single instruction, multiple data
  - MISD multiple instruction, single data
  - MIMD multiple instruction, multiple data

■ More recent (Stallings, 1993)

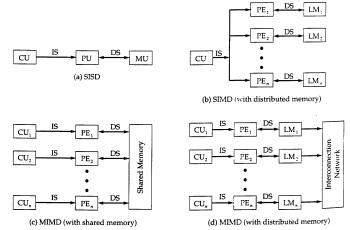
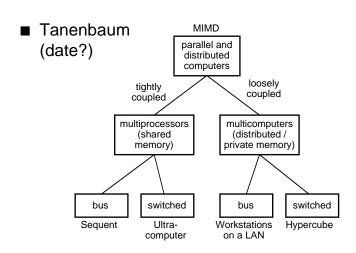


FIGURE 16.16. Alternative Computer Organizations

## Classification of MIMD Architectures



- Tightly coupled ≈ parallel processing
  - Processors share clock and memory, run one OS, communicate frequently
- Loosely coupled ≈ distributed computing
  - Each processor has its own memory, runs its own OS (?), communicates infrequently

Fall 2000, Lecture 02

### MIMD vs. SIMD (Hord 2.2)

#### ■ MIMD:

- Relatively few powerful processors
- Control level parallelism that assigns a processor to a unit of code
- Typically either distributed memory or shared memory; can have memory contention
- Needs good task scheduling for efficiency
- Each processor runs its own instruction sequence
- Each processor works on a different part of the problem
- Each processor communicates data to other parts
- Processors may have to wait for other processors or for access to data

## MIMD vs. SIMD (Hord 2.2) (cont.)

#### ■ SIMD:

- Many simple processors
- Data level parallelism that assigns a processor to a unit of data
- Typically distributed memory; can have data communication problems
- Needs good processor utilization for efficiency
- All processors are given the same instruction
- Each processor operates on different data
- Processors may "sit out" a sequence of instructions

5 Fall 2000, Lecture 02 6 Fall 2000, Lecture 02