## **Sieve of Eratosthenes**

- Find the prime numbers less than or equal to some positive integer n
  - Begin with a list of natural numbers 2, 3, 4, ..., n
  - Remove composite numbers from the list by striking multiples of 2, 3, 5, and successive primes
  - After each striking, next unmarked natural number is prime
  - Sieve terminates after multiples of largest prime less than or equal to have been struck from the list
- Sequential algorithm
  - Boolean array containing numbers being sieved,integer corresponding largest prime found so far, integer keeping track of multiples of current prime

#### A Control-Parallel Approach (cont.)

- Problems and inefficiencies:
  - Processor accesses variable containing current prime, searches for next unmarked value, then updates variable containing current prime
    - Two processors could do this at once
  - Processor could end up sieving multiples of a composite number
- How much speedup can we get?
  - Suppose n = 1000

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- Sequential algorithm
  - Multiples of 2: ((1000-4)+1)/2=997/2=498
  - Multiples of 3: ((1000–9)+1)/3=992/3=330
  - Sum = 1411 (number of "steps")
- 2 PEs gives speedup 1411/706=2.00
- 3 PEs gives speedup 1411/499=2.83
- 4 PEs is same, so upper bound is 2.83

# A Control-Parallel Approach

- Control parallelism refers to applying <u>different</u> operations to different data elements simultaneously
  - Shared-memory MIMD, Distributed-memory MIMD
- Control-parallel Sieve
  - Each processor works with a different prime, and is responsible for striking multiples of that prime and identifying a new prime number
  - Each processor starts marking...
  - Shared memory contain boolean array containing numbers being sieved, integer corresponding largest prime found so far
  - PE's local memories contain integer keeping track of multiples of its current prime (each working with different prime)

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## A Data-Parallel Approach

- Data parallelism refers to using multiple PEs to apply the <u>same</u> operation to different data elements simultaneously
  - Shared-memory MIMD, Distributed-memory MIMD, Distributed-memory SIMD
- Data-parallel Sieve

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- Each processor works with a same prime, and is responsible for striking multiples of that prime from a segment of the array of natural numbers
- Assume we have p processors, where p << sqrt(n)</li>
  - Each processor gets no more than ceiling(n/p) natural numbers
  - All primes less than sqrt(n), as well as first prime greater than sqrt(n) are in list controlled by first processor

## A Data-Parallel Approach (cont.)

- Data-parallel Sieve (cont.)
  - Algorithm

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- Processor 1 finds next prime, broadcasts it to all PEs
- Each PE goes through their part of the array, striking multiples of that prime (performing same operation)
- Continues until first processor reaches a prime greater than sqrt(n)
- How much speedup can we get?
  - Suppose n = 1,000,000
  - There are 168 primes less than 1,000, the largest of which is 997
  - Maximum execution time = (ceil(ceil(1,000,000/50)/2)+ ceil(ceil(1,000,000/50)/3)+ ceil(ceil(1,000,000/50)/5)...)etime
  - Communication time = 168(50–1)ctime

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A Data-Parallel Approach (cont.)

- How much speedup can we get? (cont.)
  - Speedup is not directly proportional to the number of PEs it's highest at 11 PEs
    - Computation time is inversely proportional to the number of processors used
    - Communication time increases linearly
    - After 11 processors, increase in communication time is higher than decrease in computation time, and total execution time increases
  - How about I/O time?
    - Have to output 78,498 primes!
    - I/O time is constant because output must be performed sequentially
    - This sequential code limits the speedup
    - Amdahl's law says that the fraction of operations that must be performed sequentially limits the maximum speedup possible (more on this later in the course)