Dealing with Deadlock (Review)

- **The Ostrich Approach** — stick your head in the sand and ignore the problem
- **Deadlock prevention** — prevent deadlock from occurring by eliminating one of the 4 deadlock conditions
- **Deadlock detection** algorithms — detect when deadlock has occurred
  - **Deadlock recovery** algorithms — break the deadlock
- **Deadlock avoidance** algorithms — consider resources currently available, resources allocated to each process, and possible future requests, and only fulfill requests that will not lead to deadlock

Resource-Allocation Graph (Review)

- The deadlock conditions can be modeled using a directed graph called a resource-allocation graph (RAG)
  - 2 kinds of nodes:
    - **Boxes** — represent resources
      - Instances of the resource are represented as dots within the box
    - **Circles** — represent processes
  - 2 kinds of (directed) edges:
    - **Request edge** — from process to resource — indicates the process has requested the resource, and is waiting to acquire it
    - **Assignment edge** — from resource instance to process — indicates the process is holding the resource instance
  - When a request is made, a request edge is added
    - When request is fulfilled, the request edge is transformed into an assignment edge
    - When process releases the resource, the assignment edge is deleted

Interpreting a RAG With Single Resource Instances (Review)

- If the graph does not contain a cycle, then no deadlock exists
- If the graph does contain a cycle, then a deadlock does exist

Deadlock Detection (Single Resource of Each Type)

- If all resources have only a single instance, deadlock can be detected by searching the resource-allocation graph for cycles
  - Silberschatz defines a simpler graph, called the wait-for graph, and searches that graph instead
    - The wait-for graph is the resource-allocation graph, minus the resources
    - An edge from p1 to p2 means p1 is waiting for a resource that p2 holds (here we don’t care which resource is involved)
- One simple algorithm:
  - Start at each node, and do a depth-first search from there
  - If a search ever comes back to a node it’s already found, then it has found a cycle
Interpreting a RAG With Multiple Resource Instances

- If the graph does not contain a cycle, then no deadlock exists.

- If the graph does contain a cycle, then a deadlock may exist.

- With multiple resource instances, a cycle is a necessary (but not sufficient) condition for deadlock.

Deadlock Detection (Multiple Resources of Each Type)

- This algorithm (Coffman, 1971) uses the following data structures:

  - **Existing Resources** (E1, E2, E3, ..., Em)
  - **Available Resources** (A1, A2, A3, ..., Am)
  - **Current Allocation**
    - C11 C12 C13 ... C1m
    - C21 C22 C23 ... C2m
    - :     :     : ...     :
    - Cn1 Cn2 Cn3 ... Cnm
  - **Request**
    - R11 R12 R13 ... R1m
    - R21 R22 R23 ... R2m
    - :     :     : ...     :
    - Rn1 Rn2 Rn3 ... Rnm

- n processes, m types of resources

  - **Existing Resources** vector tells number of resources of each type that exist.
  - **Available Resources** vector tells number of resources of each type that are available (unassigned to any process).
  - i-th row of **Current Allocation** matrix tells number of resources of each type allocated (assigned) to process i.

- Processes may have unfulfilled requests

  - i-th row of **Request** matrix tells number of resources of each type process i has requested, but not yet received.

- Notation: comparing vectors

  - If A and B are vectors, the relation A ≤ B means that each element of A is less than or equal to the corresponding element of B (i.e., A ≤ B if Ai ≤ Bi for 0 ≤ i ≤ m).
  - Furthermore, A < B if A ≤ B and A ≠ B.
Deadlock Detection Algorithm
(Multiple Resources of Each Type)

- Operation:
  - Every process is initially unmarked
  - As algorithm progresses, processes will be marked, which indicates they are able to complete, and thus are not deadlocked
  - When algorithm terminates, any unmarked processes are deadlocked

- Algorithm:
  1. Look for an unmarked process $P_i$ for which the $i$-th row of the Request matrix is less than or equal to the Available vector
  2. If such a process is found, add the $i$-th row of the Current matrix to the Available vector, mark the process, and go back to step 1
  3. If no such process exists, the algorithm terminates

Deadlock Detection Example
(Multiple Resources of Each Type)

<table>
<thead>
<tr>
<th>Existing Resources</th>
<th>Available Resources</th>
</tr>
</thead>
<tbody>
<tr>
<td>(4 2 3 1)</td>
<td>(2 1 0 0)</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Current Allocation</th>
<th>Request</th>
</tr>
</thead>
<tbody>
<tr>
<td>0 0 1 0</td>
<td>2 0 0 1</td>
</tr>
<tr>
<td>2 0 1 0</td>
<td>1 0 1 0</td>
</tr>
<tr>
<td>0 1 2 0</td>
<td>2 1 0 0</td>
</tr>
</tbody>
</table>

resources = (tape drive plotter printer CDROM)

- Whose request can be fulfilled?
  - Process 1 — no — no CDROM available
  - Process 2 — no — no printer available
  - Process 3 — yes — give it the requested resources, and after it completes and releases those resources, $A = (2 2 2 0)$
  - Process 1 still can’t run (no CDROM), but process 2 can run, giving $A = (4 2 2 1)$
  - Process 1 can run, giving $A = (4 2 3 1)$

After Deadlock Detection:
Deadlock Recovery

- How often does deadlock detection run?
  - After every resource request?
  - Less often (e.g., every hour or so, or whenever resource utilization gets low)?

- What if OS detects a deadlock?
  - Terminate a process
    - All deadlocked processes
    - One process at a time until no deadlock
      - Which one?
      - One with most resources?
      - One with less cost?
        - CPU time used, needed in future
        - Resources used, needed
        - That’s a choice similar to CPU scheduling
    - Is it acceptable to terminate process(es)?
      - May have performed a long computation
        - Not ideal, but OK to terminate it
      - Maybe have updated a file or done I/O
        - Can’t just start it over again!

- Any less drastic alternatives?
  - Preempt resources
    - One at a time until no deadlock
    - Which “victim”?
      - Again, based on cost, similar to CPU scheduling
  - Is rollback possible?
    - Preempt resources — take them away
    - Rollback — “roll” the process back to some safe state, and restart it from there
      - OS must checkpoint the process
        - Write its state to a file
        - Could roll back to beginning, or just enough to break the deadlock
      - This second time through, it has to wait for the resource
      - Has to keep multiple checkpoint files, which adds a lot of overhead
  - Avoid starvation
    - May happen if decision is based on same cost factors each time
    - Don’t keep preempting same process (i.e., set some limit)